# Purdue University Purdue e-Pubs

College of Technology Masters Theses

College of Technology Theses and Projects

4-18-2012

# Virtualization in High-Performance Computing: An Analysis of Physical and Virtual Node Performance

Glendon M. Jungels

Purdue University, gjungels@purdue.edu

Follow this and additional works at: http://docs.lib.purdue.edu/techmasters

Part of the <u>Digital Communications and Networking Commons</u>, and the <u>Other Computer Engineering Commons</u>

Jungels, Glendon M., "Virtualization in High-Performance Computing: An Analysis of Physical and Virtual Node Performance" (2012). *College of Technology Masters Theses.* Paper 64. http://docs.lib.purdue.edu/techmasters/64

This document has been made available through Purdue e-Pubs, a service of the Purdue University Libraries. Please contact epubs@purdue.edu for additional information.



Purdue

West Lafavette. Indiana



n	0	ı	0	g	V

# Performance 3: nd Virtual Node ce

uirements for the in Technology

Report

<u>Date</u>	

# Contents

College of Technology	1
Executive Summary	4
Introduction	5
Statement of the Problem	6
Significance of the Problem	7
Statement of Purpose	7
Project Background	9
Relevant History	9
Definitions	10
Review of Literature	11
Project Assumptions	
Project Delimitations	
Project Limitations	
Procedures Employed	19
File-system benchmarking	19
Data from file-system benchmarking	
Large File Copy	
Data from file-system benchmarking	24
Small File Copy	24
Data file-system benchmarking	28
Small File Writes	28
Procedures Employed	
HPC benchmarking	
Data from HPC benchmarking	
NPB - BT	
NPB - CG	
NPB - EP	35
NPB – FT	
NPB - IS	
NPB - LU	39
NPB - MG	41
NPB - SP	42

Appendices	50
References	Error! Bookmark not defined.47
Opportunities for further research	46
Conclusions, recommendations and financial implications	45
XHPL	44

#### **Executive Summary**

Advances in computing technology have changed the high-performance computing landscape. Powerful hardware, such as multiprocessor servers, and high-speed, low-latency networking technologies are available at increasingly competitive price to performance ratios. These components, combined with a modern operating system can be utilized to assemble a system capable of such actions as simulating a nuclear explosion, predicting global weather patterns, and rendering a feature length animated film.

Virtualization is a computing process that allows multiple operating systems, or multiple instances of a single operating system, to reside and function on a single computer server.

Operating hardware in this fashion offers advantages such as greater flexibility and higher utilization of the server's resources. As a result, it offers possible environmental advantages such as lower power consumption and fewer physical servers to be managed in an organization's datacenter.

The following document demonstrates the creation of a high-performance computer using modern hardware and using virtualization technology to efficiently utilize a server's computing capabilities while providing near equivalent performance as compared to equivalent physical server hardware. Resources, such as networked file-systems, and industry respected benchmarking tools are used to accumulate data from performance testing. An analysis of this testing is presented.

The results of this testing show that although creating a high-performance cluster using virtualization is possible, and offers advantages, it does not offer real-world feasibility.

Computing performance, as compared to an equivalent physical cluster, is proven to be substantially less in many of the benchmarks utilized, specifically those using high levels of

inter-node communication. Sustained file operations also frequently caused virtual servers to lock up, necessitating a reboot. The real-world implication of these results is that utilizing the resources used in this research, virtual servers should be considered for use in high performance computing where inter-node communication is minimal.

#### Introduction

Virtualization has its beginnings from the days when the IBM mainframe was the dominant computing platform. With the introduction of the PC, mainframes and mainframe-type capabilities fell to the technological wayside. Recently, the concept of virtualizing instances of operating systems has taken hold in the PC computing arena. Examples of current virtualization software in used today include EMC's VMWare, SWSoft's Virtuozzo, the open-source Xen product, Microsoft's Virtual Server and the Linux Kernel based Virtual Machine (KVM). Advancements in server processors and specialized operating system drivers allowing virtualization to take place in hardware on a server have minimized the performance penalty that occurs by virtualizing operating environments.

Previous studies have been completed to show the performance impact of utilizing non-hardware assisted virtualization techniques on high-performance clusters. The impact of such an environment using the open-source Xen software typically constitutes a performance degradation of two to five percent compared to native performance. Hardware assistance for virtualization can be used to compensate. This technology allows a more efficient operating environment in which to deploy what are commonly referred to as high-performance computing (HPC) clusters. This study hypothesizes that HPC node performance, and thus overall HPC cluster performance, can obtain near equivalent performance for CPU bound tasks compared to their native hardware counterparts.

For example, creating a 16 node HPC cluster in which each node represents a virtualized environment bound to one CPU on a 16 processor server will provide performance near equivalent to running 16 physically separate hardware nodes. This virtualized environment provides cluster administrators with potential benefits from fewer physical servers; lower power consumption, lower cooling requirements as well as easier administration for such tasks as deploying new nodes, re-claiming nodes that are not in use and more flexible node resource allocation. In a virtual server environment, resources such as memory and disk space can be assigned without physically touching hardware.

This study builds on these concepts by creating a small (4 node) virtualized cluster on a multi-processor (with hardware assistance) server in which to exercise the efficiencies mentioned above. The results of testing these efficiencies will allow for potential extrapolation of the impact on much larger cluster deployments in an IT environment in which data-center floor space is becoming ever so difficult to find and ever more expensive to build.

#### Statement of the Problem

According to Forrester's April 2006 IT Forum, in production environments the number one reason to use this technology is for flexibility. This flexibility is supported by hardware assisted virtualization, such as AMD-V (AMD 2011) and Intel VT (Intel 2011), which allows hardware acceleration of some software operations. One use for virtualization technology is to examine software impacts on clusters without deploying hardware equivalent to that found in an organization's production cluster. (Spigarolo and Davoli 2004)

Even with such hardware acceleration, virtualization does create operating overhead, particularly for I/O operations (Kiyanclar 2006). This leads to an unknown that this research seeks to address. To what degree is performance impacted when using hardware assisted

virtualization to create a virtual high-performance computing cluster?

#### Significance of the Problem

Running high-performance computing cluster nodes in a virtualized environment has several implications. One implication is that data-centers do not have to continually expand the physical space required as the data-center grows over time. There are other implications as well. Fewer, but more power efficient hardware means less power consumption and in turn less impact on the environment. This also means that fewer cooling units are needed to cool the data-center floor, allowing for more available space and requiring less consumption. Deploying less physical hardware also demands fewer technicians to keep a data-center operational each day. Collectively, these financial incentives provide compelling justification for the use of virtualization in high performance computing environments.

#### Statement of Purpose

A quick browse through any IT trade magazine will provide evidence that the world of virtualization is growing. This capability was created in the 1960's and previously built into high-end computing environments such as mainframes. Today, it is now being built into offerings that are available in the average consumer desktop PC (Huang, Liu, Abali, Panda 2006)

). Mainframe virtualization continues to grow in use and capability as well (Babcock 2007). As data-centers across the US continue to grow, virtualization takes on several important roles.

A recent study published in Information Week indicates that the majority of virtualization is taking place in order to consolidate server workloads onto fewer, but more powerful and resource abundant servers using virtualized servers. This creates several questions to consider when deploying cluster nodes as virtual servers in hardware assisted virtualized environments.

Can cluster nodes running as virtualized guests perform CPU intensive tasks as well as their

stand-alone counterpart? Previous studies indicate that using specialized hardware for Input/Output (I/O) operations can minimize performance impacts, but if the answer is no then our question must get more specific (Yu and Vetter 2008). Are there CPU performance penalties in using cluster nodes in hardware assisted virtualized environments? Does the same hold true for memory access and I/O operations such as writing to disks? What are the other advantages and limitations in deploying clusters in such a manner?

The hypothesis to be tested is that utilizing hardware assisted virtualization; high performance computing nodes can perform as nearly well as individually deployed hardware nodes when performing CPU intensive tasks. Hardware assisted virtualization allow for the Virtual Machine Monitor (VMM), the software that provides oversight to the virtualized environment, to pass on instructions previously emulated by software directly to a computer's CPU(s). (IEEE Computing Society 2005). This hypothesis becomes increasingly significant as corporate data-centers run out of physical capacity (floor space, electrical power capacity, cooling capacity). Running multiple cluster nodes on a single, multi-processor/multi-core server can aggregate equal amounts of computing power into a smaller amount of data-center space than other methods such as utilizing blade-server technology.

In order to test this hypothesis, a hardware-assisted, virtualized four node cluster will be built for the purpose of comparing computational efficiency of virtualized clusters to that of native cluster environments. The testing environment will be using the Linux operating system as it is the predominant high-performance operating system in use by clusters today (Top500.org 2011).

In short, the purpose of this study is to determine the viability of using virtualized cluster computing nodes in place of traditional nodes. In addition, the study will serve to determine if

the *potential* exists for cost and management overhead savings while retaining an equivalent level of computing performance when deploying cluster nodes as hardware-assisted virtualized servers.

#### Project Background

#### **Relevant History**

Completing searches in research databases yields few, but applicable, results at the time this initial research was completed. In 2004, instructors at the University of Bologna, Italy, constructed the Berserkr Beowulf cluster. This cluster utilizes a software virtualization (User Mode Linux) approach which allows for multiple Linux kernels to run in user-space (the area of memory in which most applications are run). The primary purpose of Berserkr is not that of performance, but rather for testing, teaching, security (of resource assignment) in a low cost environment. Specifically, this virtualized cluster is used to teach parallel programming methods in a computer science curriculum without the associated high cost of a traditional computing cluster (Spigarolo and Davoli 2004).

Faculty from Ohio State's Computer Science and Engineering department teamed up with researchers from the IBM T. J. Watson Research Center to propose that high-performance computing clusters deployed in virtualized environments have advantages over other deployments. Among the proposed advantages is ease of management, customized operating system and advanced system security by enabling only the services necessary for a program to run (a point more applicable to environments in which computing resources are shared between departments, etc). This study includes the use of the Xen virtualization technology. It acknowledges several limitations with the architecture, specifically with input/output operations. By utilizing custom software, the impacts of these limitations were minimized and enabled their

project to use a high-speed/low latency interconnects called Infiniband, rather than traditional Ethernet for communications.

Although not specifically mentioned, the timing of this study, as well as the types of hardware (no model numbers), leads one to believe that the authors did not utilize hardware assisted virtualization in creating the cluster in their study. Useful cluster performance benchmarking tools, such as NAS Parallel Benchmarks, were mentioned that will be used in this study (Huang, Liu, Abali, and Panda 2006).

#### **Definitions**

**Virtualization:** an abstraction of computing resources by hiding the physical computing resources and making it appear as a logical unit.

**Virtual Machine Monitor (VMM):** also known as a hypervisor. This is the platform that allows for multiple concurrent operating systems to run on the same physical hardware.

**Hardware Assisted Virtualization:** this is abstraction of computing resources performed at the hardware (CPU) level. This type of abstraction offers increased performance as hardware intercepts, and performs, hypervisor system calls rather than being emulated by software.

**Cluster:** in the context of this research paper, a cluster is a group of machines that work together to perform analysis of data in parallel, thus increasing the speed at which the analysis takes place.

**Parallel Processing:** the process of breaking data into pieces and spreading the analysis over a number of machines to decrease the amount of time to complete the analysis.

#### **Review of Literature**

This study represents a combination of studies in clustering, virtualization, and performance relative to running a cluster on standard hardware (not in a virtual environment) and its associated business implications. There is the engineering, or very technology specific, application of creating a cluster in a virtual environment. The research for literature needed to be representative of this view. To fulfill this research need, the three search databases that were utilized were: the Purdue University ACM Portal, Compendex and IEEE Xplorer. There is also a business, or real world implementation and benefit. Trade organizations such as Forrester and Gartner provided the necessary business outlook on the emerging technology known as virtualization, and to a lesser extent on clustering. Terms used to search these databases included: Xen, clustering, virtualization, benchmarking tools, Linux, technology, techniques, computing, hardware assisted virtualization, and various combinations of these terms. As some searches in these databases yielded few results, a search utilizing more open search tools, such as Google, was necessary to provide sufficient avenues for further research in this area.

Similar searches in the distinctly different style databases provided unique perspectives on the technologies of this study. These searches also provided much different summary areas than the search terms would have indicated. These can be categorized as follows: virtualization technologies and general impact on information technology, virtualization and utilization with clustering technology (specific to this study), and performance/benchmarking software representing a potential tool-set upon which to quantify how closely a virtualized cluster compares to clusters utilizing traditional deployments of rack mounted and blade servers.

History suggests that the success or failure of a technology is often dictated by the entities backing it, rather than the merits of simply the technology on its own. With proper

support, a technology can gather the moment to garner the attention of investment and finance entities. According to studies, virtualization is an example of this type of technology.

Forrester conducted research in late 2005 to determine adoption trends of virtualization. This research included 56 North American companies with 500 or more employees. In addition, the firm conducted a roundtable discussion at its April 2006 IT Forum. The results support virtualization technology as one that will have an ongoing impact on IT and business operations. Sixty percent of the respondents in the combined roundtable and survey reported use of some type of virtualization technology (Gillett and Schreck 2006). In many instances, virtualization technology is used for testing and development. Due to the ease of setting up a virtualized cluster to mimic a production cluster, but on less capable hardware, this strategy is also used to examine potential software impacts on high-performance computing clusters (Spigarolo and Davoli 2004).

The Forrester research also indicated that virtualization is being deployed in production environments as well for a variety of purposes such as file and prints sharing, web serving, serving custom applications, infrastructure roles such as DNS and DHCP, and a multitude of other purposes. The primary reason cited for using virtualization is flexibility, followed by consolidation and disaster recovery purposes. Other tangible and measurable benefits mentioned included floor space savings in the data-center, reduced energy consumption and reduced cooling needs.

The Forrester results are supported by The Wall Street Journal on March 6, 2007 entitled "Virtualization is Pumping Up Servers—Software that Enables Use of Fewer Machines May Cut Hardware Sales." As a demonstration, the article describes a company that has consolidated servers using virtualization technology and eliminated 134 servers, with more than three dozen to

be phased out by the end of 2007. Using technology similar to that utilized in HPC clusters, shared computing capacity was spread across the server farm using Virtual Iron From Virtual Iron Software, Inc. (Lawton and Clark 2007).

The impacts virtualization has on saving data-center floor space, energy consumption and reduced cooling are quite real. Traditional blade servers and other high density server deployments require massive amounts of power per square inch of data-center floor space. This equates to a fewer number of severs per rack in a data-center in order to maintain required electrical and cooling needs.

A study conducted by International Data Group (IDG) and Hewlett Packard concluded exactly this. Using AMD's newest sixteen core Opteron server processors, the group was able to virtualize workloads, maintain equivalent performance levels, and cut power consumption. The white paper cites being able achieve the following (IDG 2011):

- Up to 50% greater throughput in the same power and thermal footprint
- Load 33% more virtual machines per server
- Fit more servers within the existing power allotment

As industry and deployment dictate, it appears that virtualization is a legitimate technology that is being increasingly deployed in production environments today. This leads one to question if this same virtualization technology has been used in a HPC cluster.

Other sources of information specific to virtual clusters include a study at the National Center for Supercomputing Applications that focuses on creating on-demand clusters, as well as a study from Argonne National Laboratory suggesting the use of virtual clusters in support of national grids in an on-demand fashion as resources allow such allocation. Neither of these two additional studies utilized hardware assisted virtualization as the technology had not yet been

released at the time the studies were completed.

This study relies heavily on comparing performance of cluster nodes deployed as virtual servers to cluster nodes deployed as hardware servers. It is appropriate that a review of benchmarking tools and literature be completed. Unfortunately, searches in the preferred databases, such as IEEE Xplorer, did not provide suitable results and thus justified the use of a more open search engine using Google. The results returned were overwhelming in number. The research portion of the study consumed the greatest amount of time. This research did uncover a number of tools to be used in completing this study. These tools include both commercial and open-source (freely available) benchmarking tools.

The Ohio State/IBM virtualized cluster study made use of Numerical Aerodynamic Simulation (NAS) Parallel Benchmark. This package was developed by NASA at the Ames research center to test the efficiency of parallel processing systems, specifically those that are used for Computational Fluid Dynamics. NASA developed this package of benchmarks to be as generic as possible in order to provide a generic set of tools applicable across a variety of architectures (Bailey, Barszcz, Barton, Browning, Carter, Dagum, Fatoohi, Finebeg, Frederickson, Lasinski, Schreiber, Simon, Venkatakrishnan, Weeratunga 1994). NAS Parallel Benchmark tests CPU, memory, I/O, and network response and outputs the results into a commadelimited text file that can be easily imported into an Excel spreadsheet for more detailed analysis and graphical presentation.

Other tools identified on sourcforge.net include: LMBench, Procbench, Sysbench as well as commercial products such as Sisoft's SANDRA product and the SPEC products from Spec.org. The focus of this research will be on utilizing freely available, open technologies.

This literature review concludes that previous work has been completed in the area of

utilizing virtualization techniques in high-performance clusters. Because of the relative newness (released roughly in the three months) of the technology, a virtualized cluster using hardware assisted virtualization could not be located. This study will build on the excellent work from the Ohio State/IBM study and utilize a number of benchmarking tools found while completing the review, including the respected NAS Parallel Benchmark.

#### **Project Assumptions**

This study defines clustering as it applies to high-performance computing. Specifically, it will address clusters as deployed in a parallel computing environment. It will not be addressing alternate definitions of clustering, as it applies to high-availability and automatic fail-over of computing resources.

Hardware, software and underlying techniques to make them work together improve with time. Virtualization software, techniques and hardware are not exceptions to this paradigm.

Problems encountered in completing this project would likely not be encountered if completed on updated hardware and using modern virtualization software/techniques.

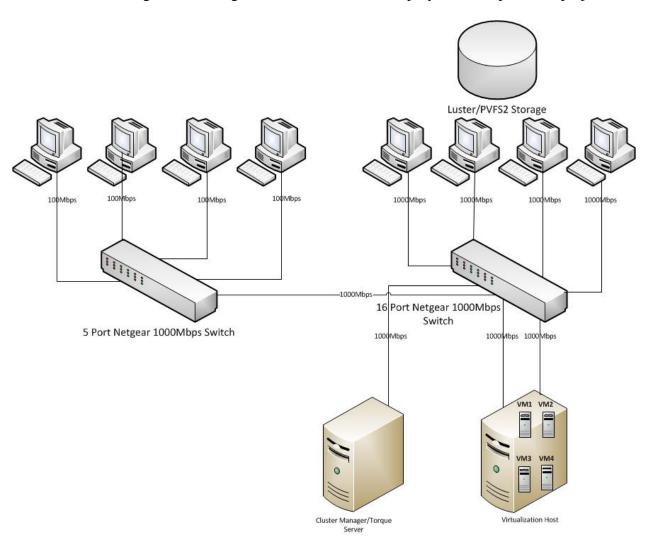
**Project Delimitations** 

The following software resources were used to complete this directed project:

Open Source Cluster	5.1.0	Cluster imaging and resource
Application Resources		tool
(OSCAR)		
Red Hat Enterprise Linux	5.0	Cluster node operating system
Citrix XenServer	6.0	Virtualization host software
Torque	2.1.8	Cluster resource manager
Maui	3.2.6p19	Cluster workload manager

Whamcloud Luster	1.8.7	Parallel file-system
Parallel Virtual FileSystem	2.0	Parallel file-system
(PVFS)		
Network File System (NFS)	3.0	Distributed file-system
Red Hat Linux Kernel	2.6.18-274	Operating system kernel
		patched to support Lustre

The following network diagram shows the network deployed to complete this project:



Based on a review of the literature, this project differs from others that are similar. Similar projects conducted utilized a single virtual processing node on each of several servers. This project utilized all virtualized processing nodes (each assigned to one processor) on a single server. This leads to differences such as resource sharing contention for access to file-systems that host the virtual machines as well as access to shared hardware such as the network controller.

This project utilized full virtualization for the virtual node rather than an alternate technique called para-virtualization that requires a modified guest kernel in order to run on the virtualization host. The earlier para-virtualization technique does not take advantage of hardware assisted virtualization thus not suitable for this study. It also requires significant kernel patching to utilize and one goal of this project was to use a kernel as close to stock as possible.

The Linux operating system used in this study utilizes a predominantly standard Linux kernel and makes use of the hardware's virtualization processing technology to operate unaltered. A single difference in the kernel was the application of a patch for the parallel file-system testing portion of this project. This was necessary to implement the Luster file-system on both the file-system side, as well as the client side. This study did not address utilizing multi-CPU virtual machine configurations.

#### **Project Limitations**

The virtualization host machine used in this study contains two processors and four processing cores. This leaves one machine at a given time sharing CPU time with the host. The physical machines (IBM desktop computers) used in the study are able to address 512 Megabytes of RAM and run at a processing speed of 1.8 GHz.

The virtualization host utilized two AMD Opteron 2212 processers, each with two cores.

Each core on the virtualization host utilizes a processing speed of 2.0 GHz. Total virtualization-host RAM of 4096 Megabytes limits each virtual machine to a maximum of less than 1024 Megabytes. Some RAM is needed by the host for processing. The version of Xen used in this project, XenServer 3.2, does not include the ability to tie a virtual host to a specific CPU. This version of XenServer also limits the number of virtual machines to 1 virtual machine per CPU, or core for multi-core machines, for a maximum of four running virtual cluster nodes.

The testing for this project consists of two associated lines of tests. The first line of testing completed benchmarks the performance of completing file-system operations on two parallel file-systems and one networked file-system by virtual cluster nodes and physical cluster nodes in independent tests. The second line of testing was completed using two well-known benchmarking programs, NAS Parallel Benchmarks and XHPL. The procedures, data, and results for the file-system testing are presented first, followed by the same format for the two benchmarking utilities.

#### **Procedures Employed**

### File-system benchmarking

The file-system testing consists of three separate tests. The first test copies a large file from the remote file-system to a location on the local node. The second copies a large number of files from a local node location to the remote file-system. The final test writes new files to the remote file-system. These tests were completed using PERL scripts, available in the appendix, on each node specific to each test, differing on each node only by the target remote file-system and folder names. Each test runs a series of 35 iterations. The time necessary to complete the iteration was logged to a text file for later analysis. The tests were completed with one, two, three and four simultaneous nodes performing the test. Only the results for four nodes are shown.

## Data from file-system benchmarking

#### **Large File Copy**

Figure 1 – Lustre Large File Copy

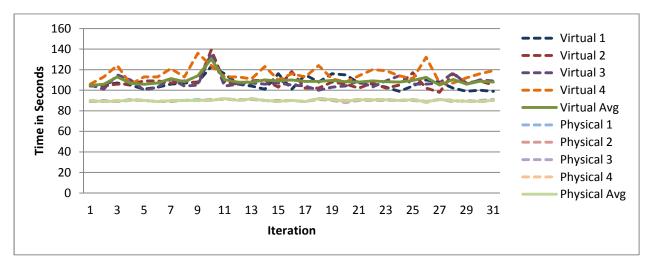


Figure 2 – PVFS Large File Copy

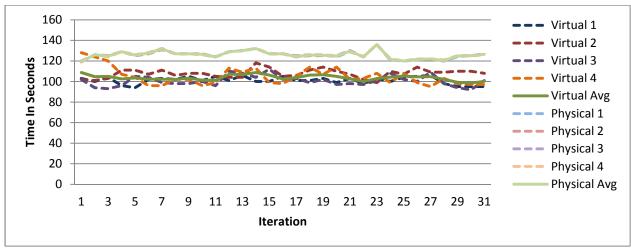


Figure 3 – NFS Large File Copy

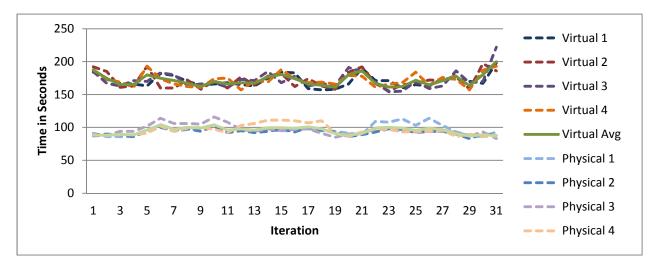
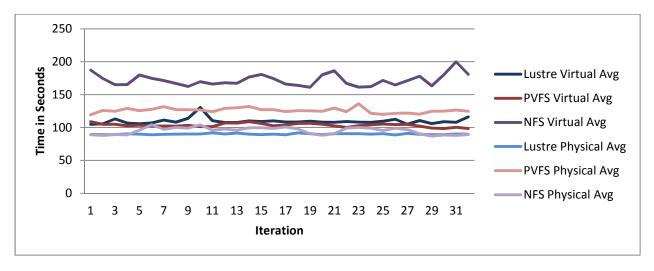


Figure 4 – Averages Large File Copy



#### Results - Large File Copy

This test was conducted for multiple purposes. Its first purpose is to compare the amount of time it takes to copy a large file of known size to a local file-system on a cluster node. Second, it compares two commonly deployed parallel file-systems used in high-performance computing clusters. Third, it compares the copy time of physical machines to that of the virtual machines. The final purpose is to compare the parallel file-systems to a commonly used conventional network file-system. The two parallel file-systems used in this test were Luster and Parallel Virtual File System version 2 (PVFS) and the single network file-system used was Network File System (NFS).

The script for this test (Appendix *Figure* 22), copies the test file from the target file-system (Lustre, PVFS, NFS) to a location on the client local file-system over a series of 35 iterations, timing and recording the copy time to a file for later analysis. This was done with a single client, two clients, three clients, and four clients to record the impacts of increasing load. The same procedure was conducted on physical cluster nodes Figures 1-3 show each file-system with four concurrent copies. Figure 1-4 displays a summary of the averages for each file-system

as accessed by both physical machines and virtual machines.

Luster was the first file-system to be tested. Figure 1 shows an immediate difference between the copy times for the virtual machines and the physical machines. The physical machine copy times show little variance. The virtual machine copy times are both longer and more erratic. There are visible peaks and valleys visible over the course of the 35 iterations that are fewer and to a lesser degree than the physical machines. There are several possible explanations for the peaks and valleys. Each physical machine has its own 100Mbps network card. This is not the case with the virtual machines as each virtual machine shares a 1000Mbps network card with another virtual machine. The 1000Mpbs connection offered by the virtualization host offers little advantage to the virtual machines as each is constrained by a 100Mbps driver. Another potential explanation for the higher copy times for the virtual machines is the simultaneous writes to the shared local storage housing the virtual machines. Other possibilities include contention for host processor cache, contention for interrupt requests and contention for disk buffers on the host local storage.

The script (Appendix *Figure* 23) copies the same file to the PVFS file-system. The results for PVFS are presented in Figure 2. Like the copies from Lustre, the copy times from PVFS vary between the physical machines and the virtual machines. The physical machines show more variation than with Lustre while the virtual machines copy times are less erratic. Unlike the copies on Lustre, the file copies on PVFS complete faster on the virtual machines than on the physical machines. With PVFS, peaks and valleys occur much more simultaneously across the machines, both physical and virtual. A closer investigation is necessary to determine the reasons PVFS and Lustre display inverse tendencies with copying large files on virtual machines versus physical machines.

The script (Appendix *Figure* 24) copies the file to the NFS file-system. Figure 3 depicts the final file-system used in this test, NFS. Although not a parallel file-system, NFS performs very well for this test. The variance pattern between the virtual machines is comparable to that of the physical machines, though the copy times are greater. The overall copy times also correspond closely to that of the parallel file-systems. For HPC workloads requiring the movement of large files from a shared file-system to local storage, NFS is as capable as both Lustre and PVFS. Years of development and use in a large variety of workloads have contributed to a stable and well performing storage file-system.

Figure 4 provides an overview of the three file-systems. The file copies completed on the physical machines did so faster than any file-system on the virtual machines. Lustre shows the best performance overall and is the best performing file-system for the virtual machines. NFS on the physical machines provides very similar copy times to Lustre, but shows the highest copy times and also the greatest variance in copy times for the virtual machines. The performance of PVFS on virtual machines is the highest while PVFS on physical machines is close to mid-way between Lustre on the physical machines and NFS on the virtual machines. These are the results for a small number of machines. As expected with each file-system, as load increased so did the times to complete an iteration of the test. Lustre, PVFS, and NFS all performed similarly on the physical machines. These results indicate that under this type of workload, parallel file-systems offer little advantage, though with loads surpassing those created by this test, the advantage would become apparent.

# Data from file-system benchmarking

## **Small File Copy**

Figure 5 - Lustre Small File Copy

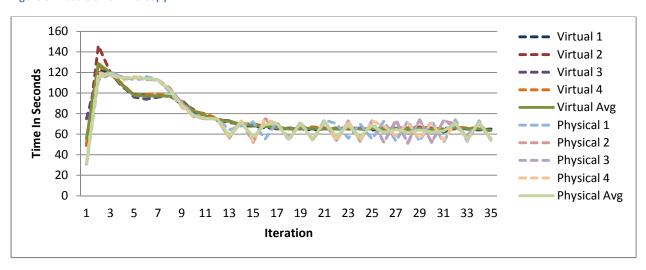


Figure 6 - PVFS Small File Copy

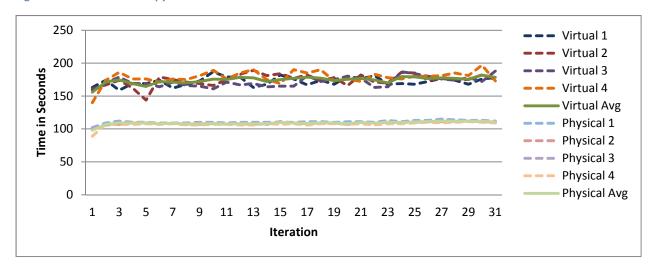
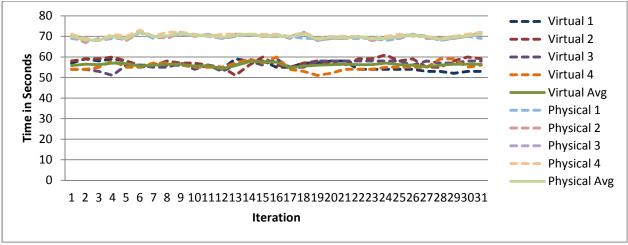
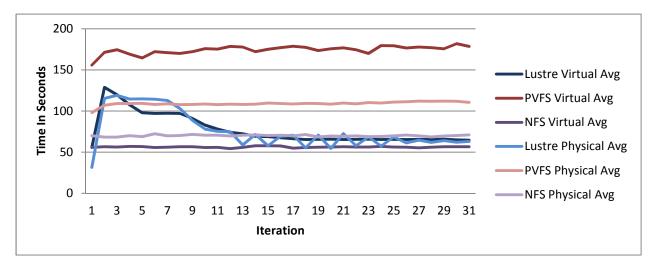


Figure 7 - NFS Small File Copy



**Figure 8 - Averages Small File Copy** 



#### Results - Small File Copy

The small file copy test copies 1000 one kilobyte files to a folder on the target file-system for 35 iterations. Each machine copies to a separate folder and the time it takes to complete the iteration is written to a file on the local machine file-system. This test was conducted to compare the copy times of virtual machines to that of the physical machines for target file-systems on Lustre, PVFS and NFS. Like the large file copy described earlier, this test was completed with one client, two clients, three clients and four clients running simultaneously. The results of the

tests are presented in Figures 5 - 8.

The script in Appendix *Figure* 25 performs the small file copy test for Lustre. Figure 5 show the results of the small file copy test on the Lustre file-system. The graph depicts some unusual results at the beginning that taper off approximately one-third of the way through the thirty five iterations of the test. Unlike the previous test, copying files to the target places continuous load on the meta-data server for the parallel file-system as new files are added. Near iteration 11, the copies stabilize and this continues until the final iteration. This is observed by tests on both the physical machines and the virtual machines. Unlike previous tests, the physical machines show variability in copy times resulting in the saw-blade look of the graph. The results also show that the copy times are almost identical within the respective groups. At various points, most notably for the virtual machines, the lines appear very close to being a single line, even with the jagged pattern of the physical machines.

Appendix *Figure* 26 is the script used to perform the small file copy test to PVFS. The results of the small file copy test for PVFS are shown in Figure 6. With PVFS, the graphs again look closer to that of the large file copy. Also once again, the virtual machines show increased variability that is not present with the physical machines. With PVFS, it is the physical machines that appear to be a single solid line. Copies on the physical machines also complete faster, unlike the large file copy in which the virtual machines perform better. One key difference between the copies on Lustre and those on PVFS is the time it to complete each iteration. The copies on PVFS take considerably longer for both the physical machines and the virtual machines.

Appendix *Figure* 27 is the script used to perform the small file copy test to NFS. Figure 7 shows the results of the small file copy test on NFS. Like PVFS, NFS shows some degree of variability in this test. This is most visible with the virtual machines, though also present with

the physical machines as well. Unlike the previous large file test with NFS, the virtual machines complete these tests faster than the physical machines. A final observation is that the slowest copies by the physical machines on NFS are as fast, or slightly faster, than the fastest copies on Lustre and the fastest virtual machine copy on NFS is nearly half the time of the fastest copy on PVFS. Like the previous large copy test, the advantages of using a parallel file-system are not realized by these tests with a small number of clients.

Figure 8 presents a summary of the three file-systems for small file copies on both virtual machines and physical machines. NFS on the virtual machines proves to be the best combination running this test, followed closely by Lustre on physical machines. Once again, the unusual beginning of the tests for Lustre on both the virtual machines and the physical machines is visible. As Lustre stabilizes, the times become an almost three way tie between Lustre on physical machines, Lustre on virtual machines, and NFS on physical machines. Like previous tests, PVFS does not excel in this test and shows the two overall highest copy times. With limited load, the simplicity of NFS proves again that it is a capable file-system.

# Data file-system benchmarking

#### **Small File Writes**

Figure 9 - Lustre Small File Write

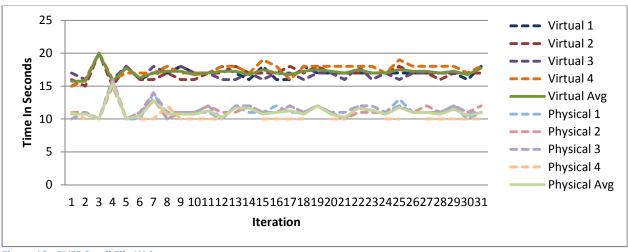
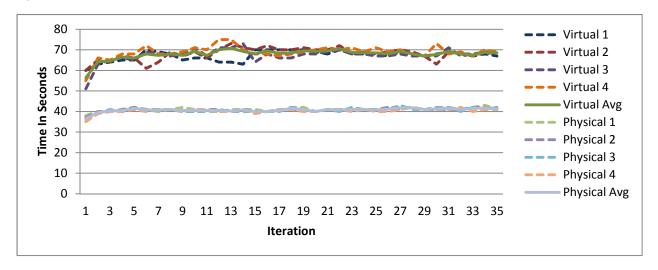


Figure 10 - PVFS Small File Write



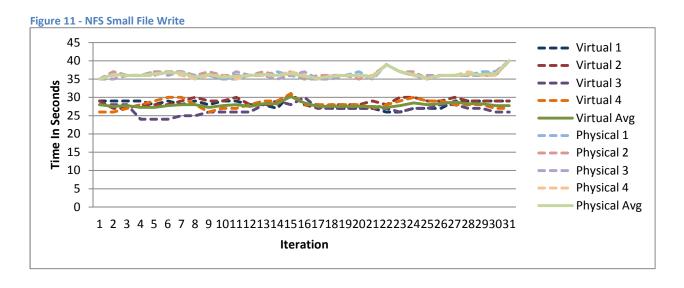
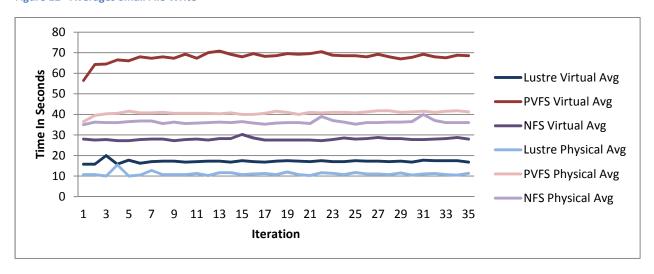


Figure 12 - Averages Small File Write



#### Results - Small File Writes

The last set of tests performed consists of writing one kilobyte files 1024 times in a single iteration and repeating this for 35 iterations. This test results in a total of thirty five megabytes written to the target file-system. The iteration was timed and the result logged to a text file on the local file-system. This test was conducted to compare how quickly a physical or virtual machine could write to a remote file-system. Operations such as this are common in high-performance computing calculations where the results are then aggregated into a final format as a

result of the processing. Analyzing the write times serves as an indicator of the speed in which a virtual machine can write a file in comparison to that of a virtual machine as well as serve as an indicator of which of the file-systems tested is the most efficient at writing small files.

The script in Appendix Figure 28 was used to complete the small file write test on Lustre. Figure 9 shows the results of the small file write test on the Lustre file-system. The first several iterations show an unusual pattern, similar that that which was observed in the small file copy test for Lustre previously discussed. After this small number of iterations, the write times even off and remain relatively consistent for the duration of the testing. Both the physical machines as well as the virtual machines display a small bit of variability, though not exaggerated like in other tests. Previous test have shown that the physical machines complete the tests faster and this is the case with this test as well. There are visible gaps in the graph. During these iterations, present only in the data for the physical machines, the machine was unable to write the file to the Lustre file-system. During those same times, other machines continued to write uninterrupted, but with higher write times than when this event is not occurring. Rather than being unable to write files for periods of time, the virtual machines display a different observable behavior. Flat lines on the graph indicate that those iterations took the exact number of seconds for those iterations. The timer does utilize whole seconds as the unit of measure, but this behavior was not observed in other tests.

Appendix *Figure* 29 is the script used to complete the small file write test on PVFS. The results for PVFS for writing small files are shown in Figure 10. The writes for PVFS show a kind of "ramp up" behavior for both the physical machines for iterations one through three. From iteration three forward, the write times are much more consistent with the physical machines again completing the iterations more quickly. The physical machines write to the

PVFS file-system with times very similar between each machine. Like previous tests, the virtual machines show more variable times between iterations, with one a section of approximately ten iterations where the times were very close. There is no missing data for the writes to PVFS as was observed in the file writes to Lustre.

The script in Appendix *Figure* 30 was utilized to complete the small file write test on NFS. Figure 11 displays the results for small file write test to the NFS file-system. Like the previous test on PVFS, both the physical machines and virtual machines were able to complete the writes for all iterations. NFS does display write behavior similar to that of Luster, showing write times of consistently the same time. Unlike Lustre, this is observed for multiple virtual machines, at times concurrently. Using NFS as the target file-system, the virtual machines also recorded lower write times than the physical machines. Write times on NFS were greater than on Luster, but lower than write times on PVFS. Though more prevalent with the virtual machines, variability in write times is minimal and there primarily in the first eleven iterations.

Figure 12 shows the summary results for small file writes for all three file-systems from both the physical and virtual machines. Luster shows the lowest write times per iteration for both physical machines and virtual machines, but as indicated previously, also had write failures from a physical machine during testing while others were able to continue. Writes to NFS from the physical machines and virtual machines show the next best write times per iteration, with all iterations completed without a failure to complete a write. PVFS on physical machines and virtual machines display the slowest write times of the file-systems tested. Like NFS, PVFS completed all iterations without a failure to write, though with a higher write time.

#### **Procedures Employed**

#### **HPC** benchmarking

As with the file-systems tests, scripting, as shown in Appendix Figures 32 - 34, was used to automate running the benchmark. These scripts controlled the timing and submission of the individual benchmarks to the cluster resource manager for assignment and completion. Upon completion of all HPC benchmarks, a PERL script was run against the output files to consolidate the results into a single text file that was later imported into Excel for analysis.

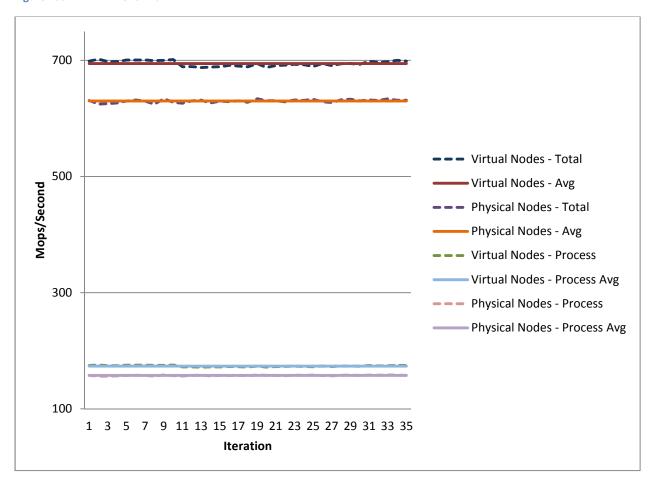
#### Benchmark Description - NAS Parallel Benchmarks

NASA Advanced Supercomputing (NAS) is responsible for the creation of the NAS Parallel Benchmarks. This small set of parallel applications was written at NASA Ames Supercomputing Center as a way to benchmark new high-performance computers being deployed. These applications utilize Computational Fluid Dynamics (CFD) equations commonly in use as part of NASA research. Together, the applications provide a generalization of the performance a new supercomputer can be expected to achieve when applied to real-world problems. The eight benchmarks test a variety of characteristics including memory access, node to node communication and processor performance. Each benchmark has a number of classes that can be utilized. The classes differ in the problem size utilized. For classes A-C, problem sizes increase by four times over the previous class. Classes D, E and F, used for testing very large supercomputers, utilize a step of sixteen times over the previous step. The W class is present, but now deprecated and the S class is intended to provide a quick test of functionality. This research utilized class B for all benchmarks presented. This problem size kept the cluster working longer than the class A benchmark, but did not exceed the memory capability of the hardware utilized as with the class C benchmark

# **Data from HPC benchmarking**

**NPB-BT** 

Figure 133 - NPB BT Benchmark



#### Results - NPB BT

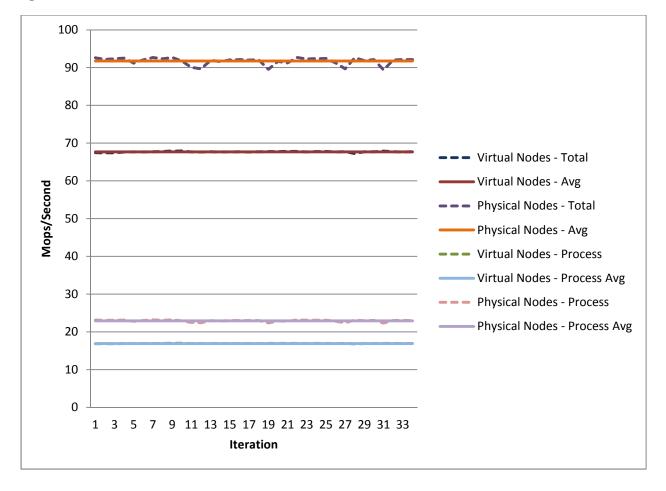
Block Tri-diagonal solver. (NASA 2012)

Figure 13 shows the results of the BT benchmark for both the physical cluster nodes as well as the virtual cluster nodes. In this benchmark the virtual cluster nodes were able to establish a slight advantage over the physical cluster nodes. A closer inspection of the lines reveals that for the algorithm used in this benchmark, the data shows less variation for the

physical cluster nodes. Although more defined in the data for the virtual nodes due to a slight drop in processing around iteration 10, both sets of cluster nodes show a slight upward rise in processing as the iterations progress. This could indicate a processor caching of frequently used data in the benchmark.

**NPB - CG** 

Figure 14 - NPB CG Benchmark



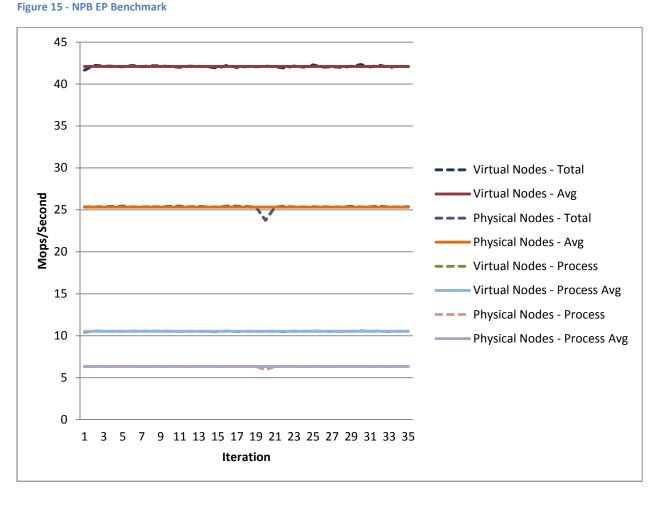
#### Results - NPB CG

The CG benchmark is a conjugate gradient method used to compute an approximation to the smallest eigenvalue of a large, sparse, symmetric positive definite matrix. This kernel is typical of unstructured grid computations in that it tests irregular long distance communication, employing unstructured matrix vector multiplication. (Bailey, Barszcz, Barton 1994)

Figure 14 show the results of the CG benchmark. The physical cluster nodes are able to produce a higher benchmarks score. The irregular communication present itself in the graph via a rolling wave shape in the data. This shape is also present in graph for the virtual cluster nodes, though less prevalent. The virtual cluster nodes display less variability in the data as the points are closer to the test average. The data for the physical cluster nodes is the same but with more defined valleys where processing drops. This may be explained by the slower network connection present on the physical nodes.

# Etaura 4E NDD ED Davidson of

NPB - EP



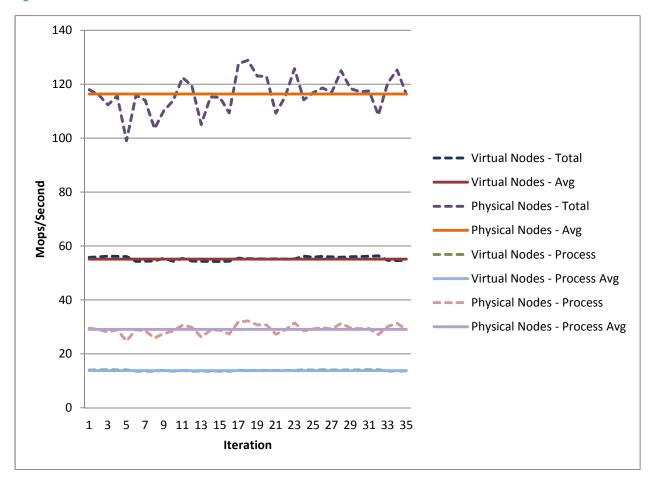
#### Results - NPB EP

The EP benchmark is the "embarrassingly parallel" kernel. It provides an estimate of the upper achievable limits for floating point performance, i.e. the performance without significant inter-processor communication (Bailey, Barszcz, Barton 1994). This benchmark is named this because there is no inter-node communication required, thus parallel to the number of cluster nodes.

Each virtual cluster node has a 200MHz processor advantage over the physical cluster nodes. This gives the virtual cluster nodes a significant advantage where there is little dependence on inter-node communication. The results are presented in Figure 15. Across thirty-five iterations of this benchmark, the virtual cluster nodes performance nearly doubles that of the physical nodes. For both node types, performance is stable, with little variation from the overall average. The per-process performance is almost indistinguishable from the per-process average.

#### NPB - FT

Figure 16 - NPB FT Benchmark



#### Results - NPB FT

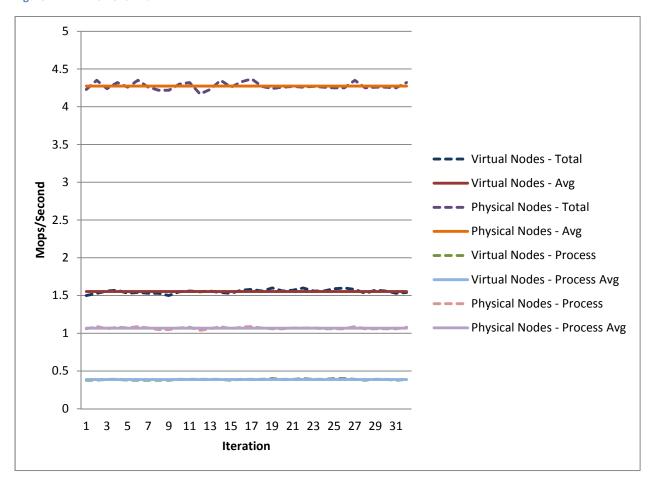
The FT benchmark is a 3-D partial differential equation solution using Fast Fourier Transforms. This kernel performs the essence of many "spectral" codes. It is a rigorous test of long-distance communication performance (Bailey, Barszcz, Barton 1994).

The results of this benchmark are presented in Figure 16. The physical cluster nodes are able to out-perform the virtual cluster nodes. The impact of rigorously testing the long distance communication performance is evident. The physical cluster nodes show a great degree of variability over the entire thirty five iterations. The performance of the virtual cluster nodes

appears to be less impacted in terms of variability, though the performance is roughly half overall.

#### **NPB-IS**

Figure 17 - NPB IS Benchmark



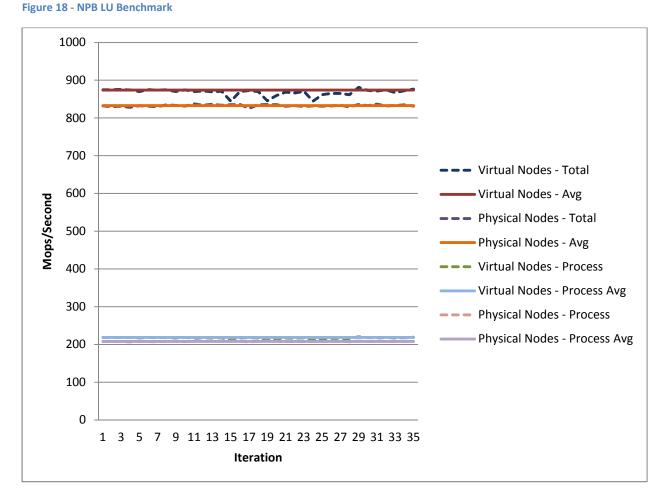
#### Results - NPB IS

The IS benchmark is a large integer sort. This kernel performs a sorting operation that is important in "particle method" codes. It tests both integer computation speed and communication performance (Bailey, Barszcz, Barton 1994).

Figure 17 shows the results of the IS benchmark. A processor speed advantage by the virtual cluster nodes does not provide an edge as in the embarrassingly parallel benchmark. This

benchmark does use node to node network communication. This is an area of weakness of the virtual cluster nodes. The resulting performance shown is a performance advantage of almost three times by the physical cluster nodes. Both types of nodes display some variability with the physical nodes showing slightly more over the course of thirty iterations. It is possible that this can be attributed to their slower 100Mbps network connection as the virtual nodes also display this behavior, but with smaller peaks and valleys.

NPB - LU



### Results - NPB LU

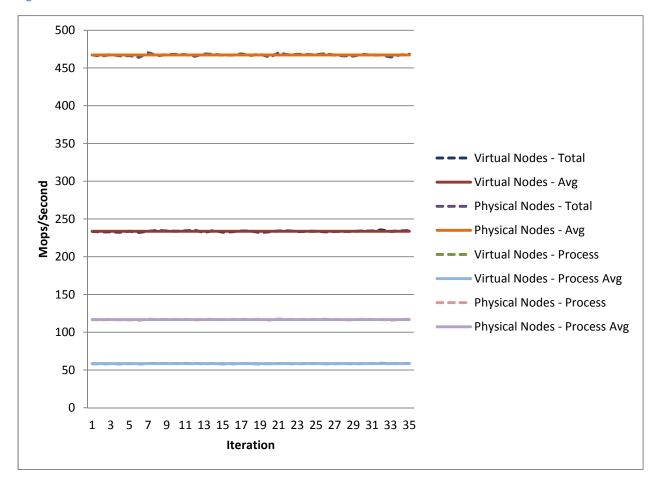
The LU benchmark solves a synthetic system of non-linear partial differential equations

using a Lower-Upper symmetric Gauss-Seidel kernel (NASA 2012).

Figure 18 illustrates the results of the LU benchmark. This was the highest performing benchmark across those utilized in this testing. The virtual cluster nodes were able to obtain nearly 1000 MOPS/Second, followed closely by the physical cluster nodes. The data for the physical nodes closely follow the average while the virtual cluster nodes vary across several iterations. This data shows the performance impact across the virtual cluster for an event that was likely an even happening on the virtualization host itself rather than an even in the benchmark itself. The same pattern of performance drop is not visible on the physical hosts at any point during the thirty five iterations of testing.

#### **NPB-MG**

Figure 19 - NPB MG Benchmark



#### Results - NPB MG

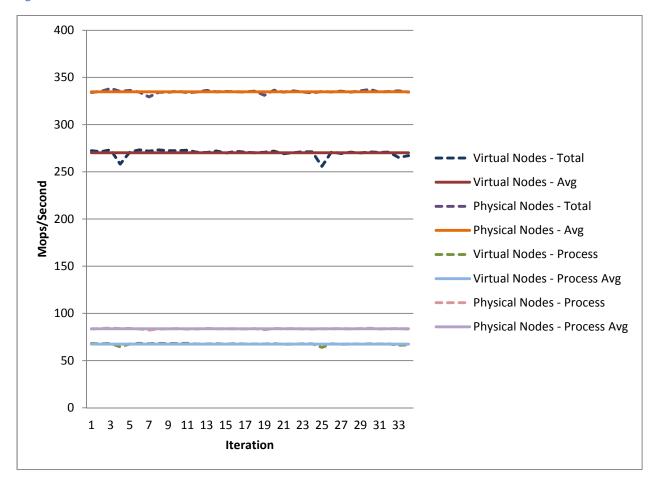
The MG benchmark is Multi-Grid on a sequence of meshes. This benchmark requires high structured long distance communication and tests both short and long distance data communication (Bailey, Barszcz, Barton 1994).

Figure 19 shows that the physical cluster nodes are able to perform better on this benchmark by nearly a factor of two. Performance aside, the virtual cluster nodes and physical cluster nodes data pattern is very similar. There is little variability from the average for either set of nodes over any of the iterations. There are no discernible peaks and valleys as seen in other

NAS benchmarks. The highly structured communication pattern is shown in the presentation of the data.

# NPB - SP

Figure 20 - NPB SP Benchmark



#### Results - NPB SP

The SP benchmark is a Scalar Pentadiagonal symmetric successive over-relaxation solver kernel for nonlinear partial differential equations (NASA 2012).

Figure 20 illustrates the results of the SP benchmark. The physical cluster nodes perform better than the virtual cluster nodes and show less variability overall. Three drops in performance pull the average down slightly. These events do not appear to be outlier data as

similar events can be seen in the data on the physical nodes. On the virtual cluster nodes, this could be an even on the virtual host itself, but as these drops in performance can be seen in the data on the physical nodes, it is likely to be a common point in the benchmark, thus impacting both groups.

Figure one through eight displays the results of running the NAS Parallel Benchmarks on the physical and virtual clusters. The highest performance on five of the eight benchmarks was obtained by the physical cluster nodes. Three of these five benchmarks specifically utilize node to node (network) communication to a larger extent. Simultaneous access to the host network resources appears to be an area where virtualization can use improvement. The two highest performing benchmarks, LU and BT, show that access to a higher clock rate processor provides an advantage in operations where little communication is necessary between the cluster nodes. In these two benchmarks, the virtual cluster nodes outperform the physical cluster nodes.

# Benchmark Description – High-Performance Linpack (XHPL)

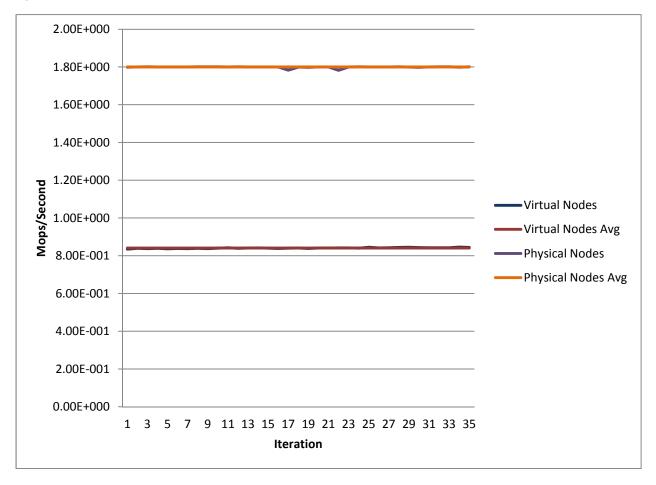
"The Linpack Benchmark is a measure of a computer's floating-point rate of execution. It is determined by running a computer program that solves a dense system of linear equations."

(Top500.org)

Unlike the NAS Parallel Benchmarks, XHPL is configurable in order to obtain maximum performance for a given parallel computer as well as to troubleshoot problem areas for new parallel computer installations. Configuration is done via a file named HPL.dat by default. The configuration options used to complete this benchmark can be viewed in Appendix *Figure 31*.

#### XHPL.

Figure 21 - XHPL Benchmark



# Results - XHPL

Figure 21 displays the results of the XHPL benchmark completed for thirty five iterations. For both the physical cluster nodes and the virtual cluster nodes, the data values show little variation. Both the physical and virtual cluster nodes show a small number of iterations that fall outside the average. The physical virtual nodes show two iterations in a short time-frame that drop below the average. The virtual nodes display the opposite. Three iterations appear above the average, while the single initial iteration falls below. The final approximately eight iterations appear to be above the iteration average as well.

### Conclusions, recommendations and financial implications

Despite the potential of utilizing high performance computer nodes in a virtualized environment, this project has uncovered a number of unexpected drawbacks. The largest of these is the negative impact on performance of simultaneous access to shared resources such as local disks, shared network adapters and parallel/networked file-systems.

The results of the NAS Parallel Benchmarks indicate that I/O is a problem for a virtualized cluster. This is shown in the figures for the CG (p. 34), FT (p. 37), IS (p. 38), MD (p. 41) and SP (p. 42) benchmarks and proceeding results analysis. The NAS Parallel Benchmarks requiring the greatest amount of inter-node communication were the worst performing. For those benchmarks requiring little inter-node communication, the virtual cluster was able to outperform the physical cluster. The figures for the BT (p. 33), EP (p. 35), and LU (p. 39) benchmarks illustrate this type of performance.

The results of the XHPL benchmark (p. 44) indicate that there is also a disparity between the capabilities of floating point operations on physical cluster nodes versus virtual cluster nodes. A goal of this project was to compare the virtual and physical cluster nodes using the same configuration. Despite having a higher CPU clock rate, the virtual cluster nodes were outperformed by a wide margin using the same configuration file for testing each. With this benchmark, physical cluster nodes produced the highest benchmark scores.

Due to the performance observed in this research, virtualized clusters appear to be a viable option for use in high-performance computing using the components specified for computing with little inter-node communication. Performance characteristics for computations requiring significant inter-node communication should be evaluated carefully prior to deployment in a virtualized environment.

Virtualized clusters do represent a potential flexibility that is more difficult to obtain with physical machines. One such opportunity to utilize this flexibility would be to use virtualized cluster resources on servers during times where utilization is low, such as off-hours and on weekends. In circumstances where top performance is the key driver, dedicated physical cluster nodes are the best solution. In situations where performance can be sacrificed for the option to multi-purpose, virtualized cluster nodes may prove to be an option.

# Opportunities for further research

There are a multitude of changes that could be made to this project for further research. Each benchmark could be tuned for optimum performance rather than focusing on maintaining consistent configuration. Another such change that would make an immediate impact on both the performance and stability of the virtual machines would be to deploy them to a dedicated SAN to eliminate local host disk/file-system issues and to allow the virtual machines to be backed up via SAN snapshots. A second possibility would be to conduct the same benchmarks again using alternate virtualization software on the host machine, such as VMWare's ESX server or the new native Linux kernel implementation KVM. A final variation would be to utilize a high-speed, low latency, interconnect such as Infiniband in order to off-set the performance penalty in the network I/O area of virtualization. This would benefit inter-node communication as well as communication with parallel file-systems should one be used.

#### References

Babcock, Charles. (2007). Virtualization, Original Mainframe Flavor. Information Week.

February 22, 2007 edition. P 22.

Bookman, Charles. (2003). Linux Clustering: Building and Maintaining Linux Clusters.

Indianapolis, IN: New Riders.

International Data Group. (2011, September). *The Power of Virtualization*. Retrieved March 18, 2012 from

http://www.serverwatch.com/ebooks/37619110/95990/1614220/119068?BLUID=201203250630 19465154662

IEEE Computing Society. (May, 2005). Virtual Machine Monitors: Current Technology and Future Trends. Retrieved Feb 14, 2007 from

http://csdl2.computer.org/persagen/DLAbsToc.jsp?resourcePath=/dl/mags/co/&toc=comp/mags/co/2005/05/r5toc.xml&DOI=10.1109/MC.2005.176.

Kiyanclar, Nadir. (2006). A Virtualization Based Solution for Secure On-Demand Cluster

Computing. Unpublished master's thesis, University of Illinois at Urbana-Champaign.

USENIX Annual Technical Conference. (2005). *Measuring CPU Overhead for I/O Processing in the Xen Virtual Machine Monitor.* Conference proceedings of the annual conference.

Vrenios, Alex. (2002). Linux Cluster Architecture. Indianapolis, IN: Sams.

http://www.vmware.com

http://www.xensource.com

http://www.virtualiron.com

Cluster File Systems, Inc. (2007). Lustre 1.6 Operations Manual. Boulder, CO: Cluster File

Systems, Inc. Retrieved on 06/13/2007.

Yu, Weikuan & Noronha, Ranjit & Liang, Shuang, & Panda, Dhabaleswar. (2006). Benefits of High Speed Interconnects to Cluster File Systems A Case Study with Lustre. Ohio State University, Department of Computer Science and Engineering.

Huang, Liu, Abali, Panda. (2006). A Case for High Performance Computing with Virtual Machines. Ohio State University, Computer Science and Engineering Department and IBM TJ Watson Research Center.

Yu, Weikuan & Vetter, Jeffery. (2008). *Xen-Based HPC: A Parallel I/O Perspective*. Eighth IEEE International Symposium on Cluster Computing and the Grid. Computer Science and Mathematics, Oak Ridge National Laboratories.

Emenecker, Wesley & Jackson, Dave & Butikofer, Joshua & Stanzione, Dan. (2006). *Dynamic Virtual Clustering with Xen and Moab*. Arizona State, Fulton High Performance Computing Institute. Retrieved on 9/15/2007 from

http://www.springerlink.com/content/uk18j8v37m24824u/

Cluster File Systems, Inc. (2007). Lustre-Datasheet.

PVFS2 Development Team. (2003). A Quick Start Guide to PVFS2. Parallel Architecture Research Laboratory, Clemson University. Retrieved on 09/15/2007 from http://www.parl.clemson.edu/pvfs/files.html.

PVFS2 Development Team. (2003). Parallel Virtual File System, Version 2. Parallel Architecture Research Laboratory, Clemson University. Retrieved on 09/15/2007 from <a href="http://www.parl.clemson.edu/pvfs/files.html">http://www.parl.clemson.edu/pvfs/files.html</a>.

Micaela Spigarolo and Renzo Davoli. 2004. Berserkr: a virtual beowulf cluster for fast

prototyping and teaching. In <em>Proceedings of the 1st conference on Computing frontiers</em> (CF '04). ACM, New York, NY, USA, 294-301. DOI=10.1145/977091.977133 http://doi.acm.org/10.1145/977091.977133

NASA Advanced Supercomputing Division. (n.d.) 2012 NAS Parallel Benchmarks. Retrieved March 1, 2012 from <a href="http://www.nas.nasa.gov/publications/npb.html">http://www.nas.nasa.gov/publications/npb.html</a>

http://www.clusterfs.com/

http://www.clustermonkey.net

http://cpan.perl.org

http://sites.amd.com/us/business/it-solutions/virtualization/Pages/server.aspx

http://www.intel.com/technology/virtualization/technology.htm?wapkw=(Intel+VT)

# Appendix A

Figure A1. Lustre Large File Copy

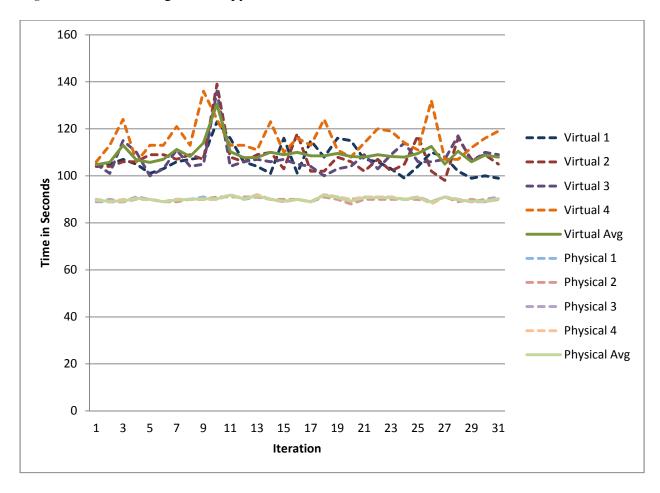


Figure A2. PVFS Large File Copy

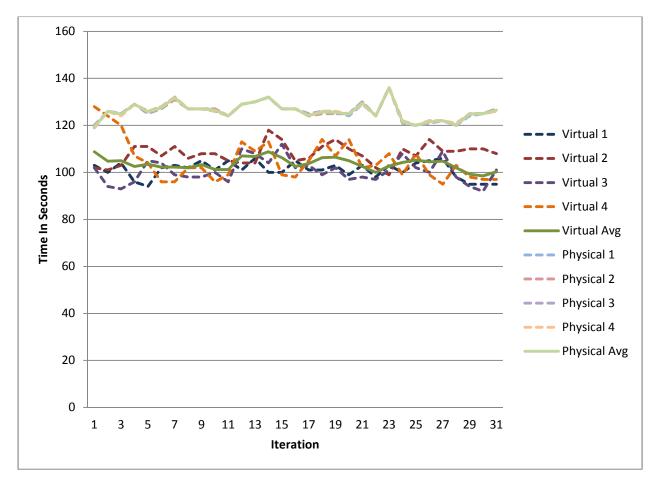


Figure A3. NFS Large File Copy

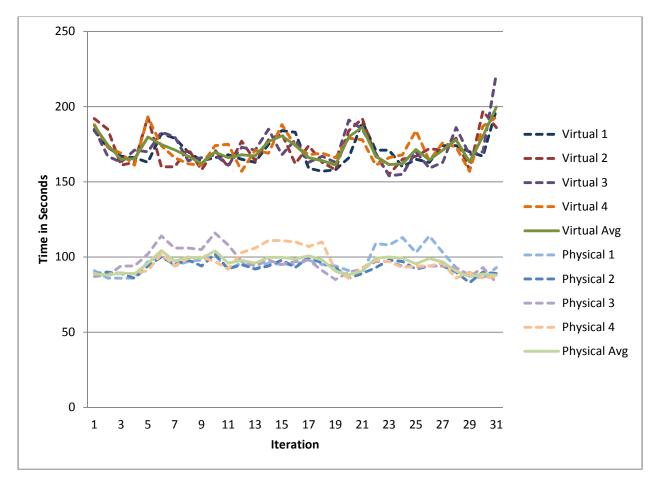


Figure A4. Averages Large File Copy

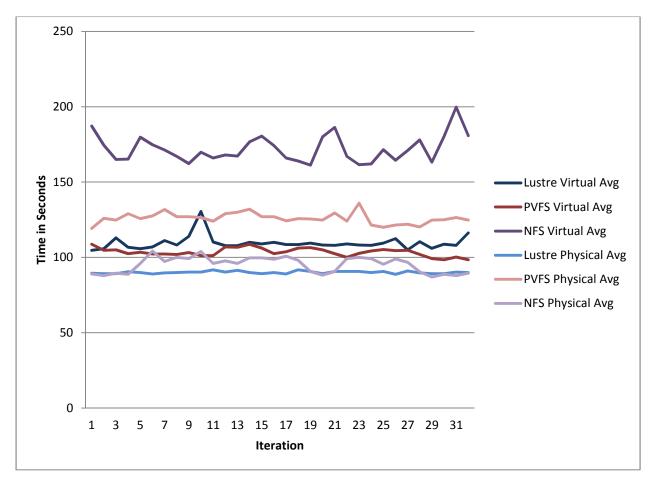


Figure A5. Lustre Small File Copy

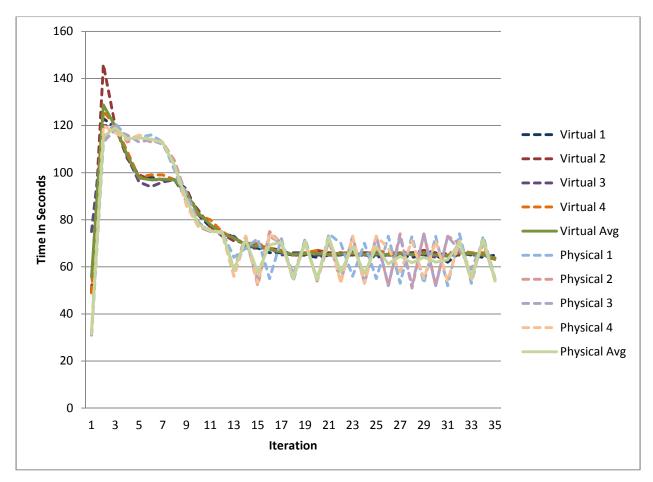


Figure A6. PVFS Small File Copy

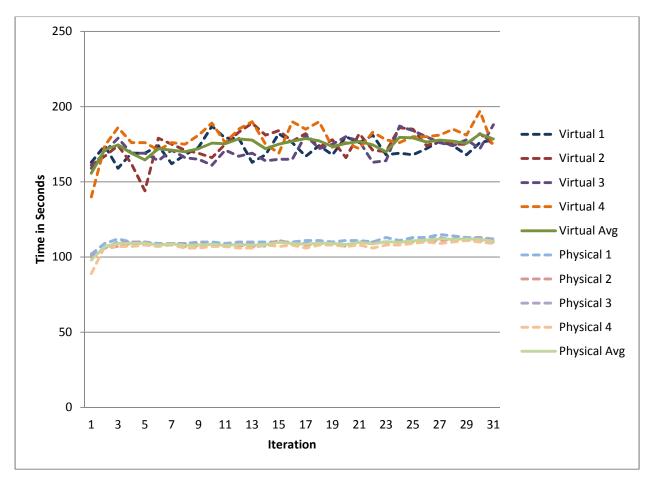


Figure A7. NFS Small File Copy

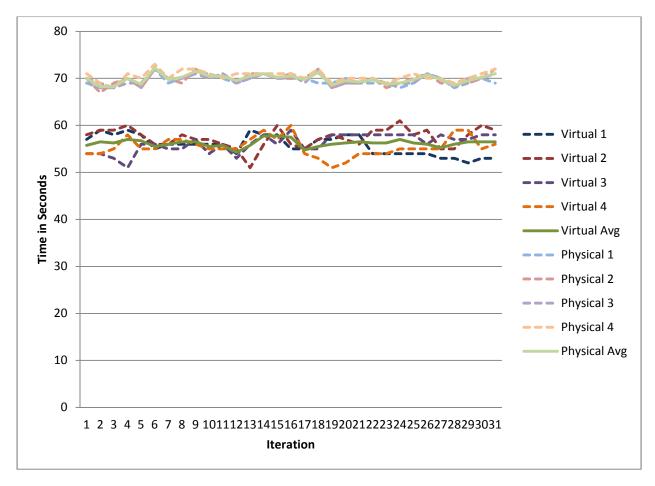


Figure A8. Averages Small File Copy

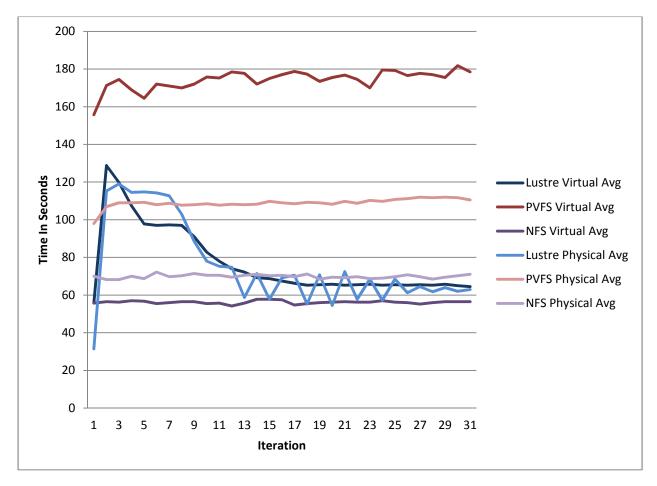


Figure A9. Lustre Small File Writes

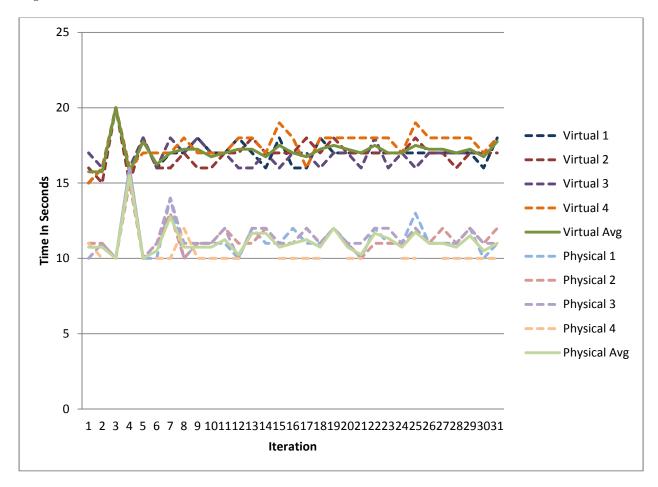


Figure A10. PVFS Small File Writes

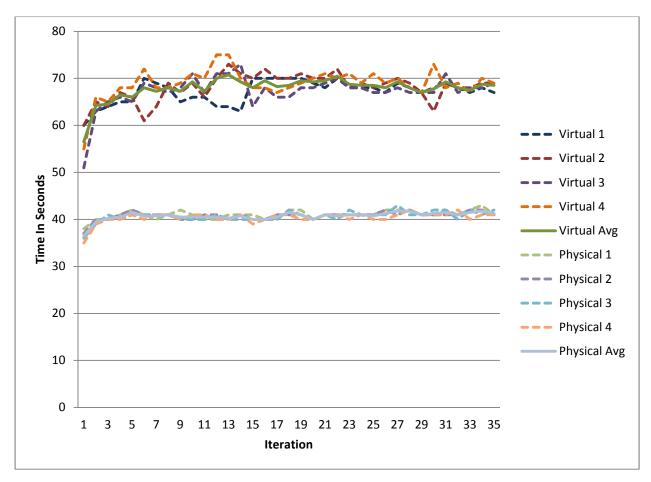


Figure A11. NFS Small File Writes

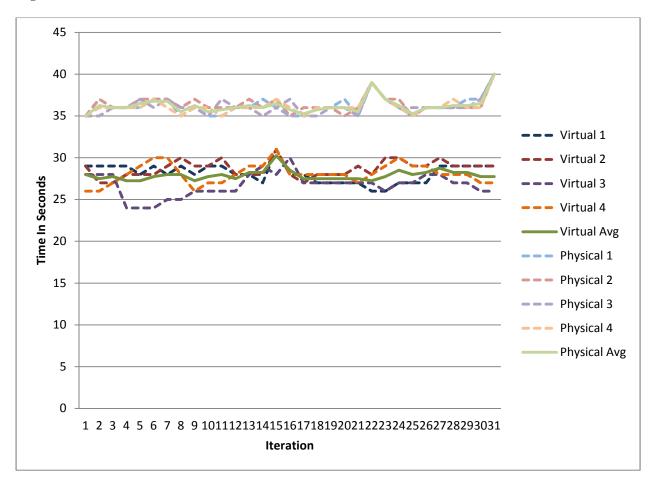


Figure A12. Averages Small File Writes

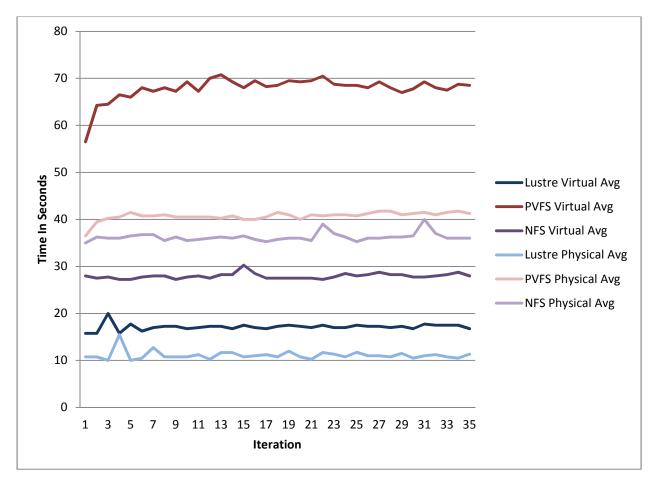


Figure A13. NAS Parallel Benchmarks BT

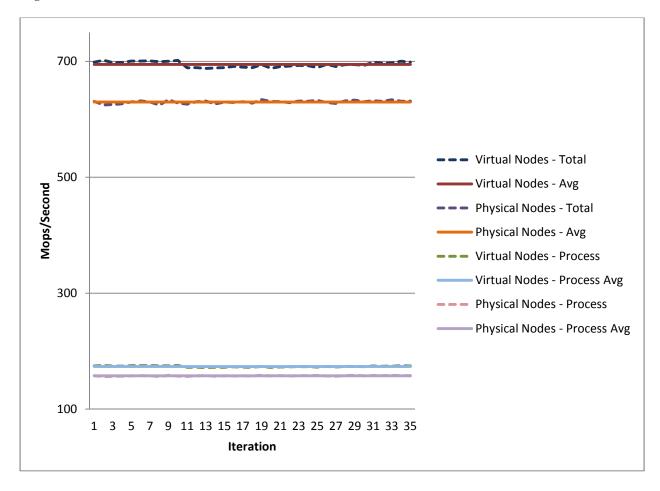


Figure A14. NAS Parallel Benchmarks CG

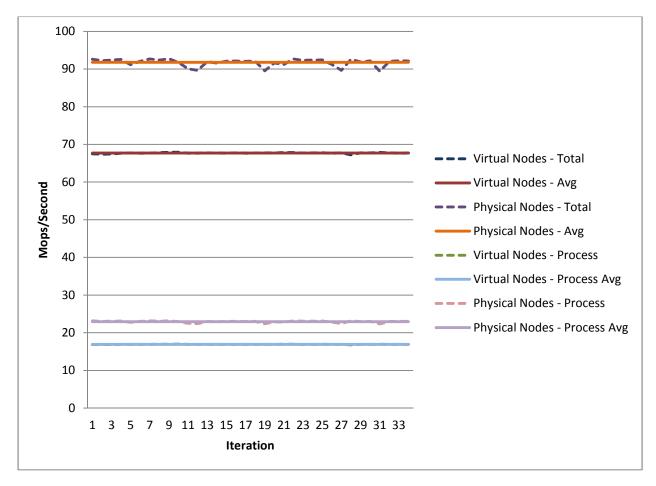


Figure A15. NAS Parallel Benchmarks EP

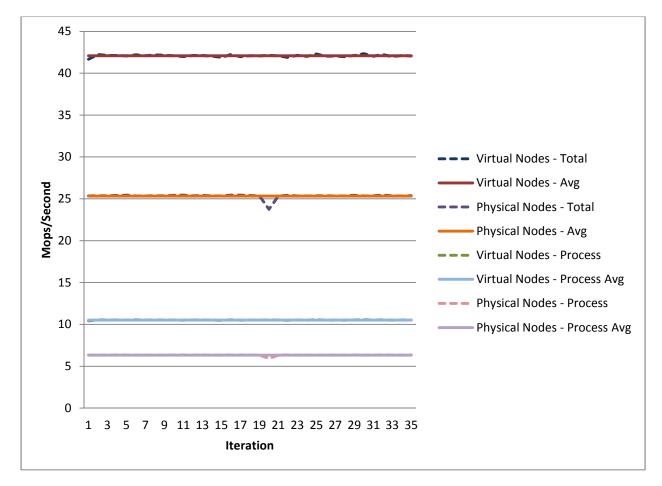


Figure A16. NAS Parallel Benchmarks FT

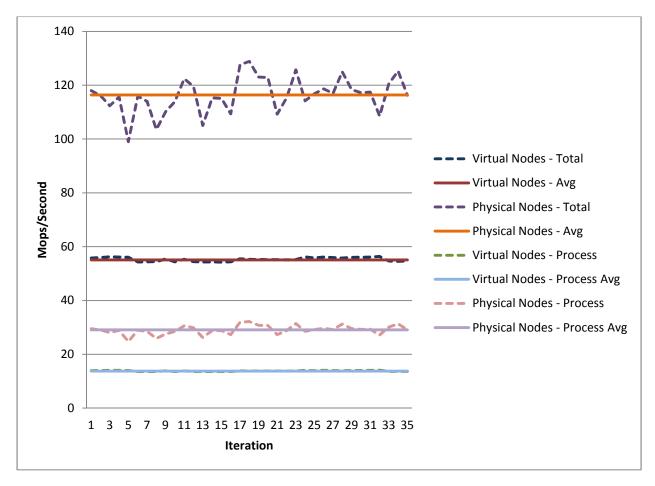


Figure A17. NAS Parallel Benchmarks IS

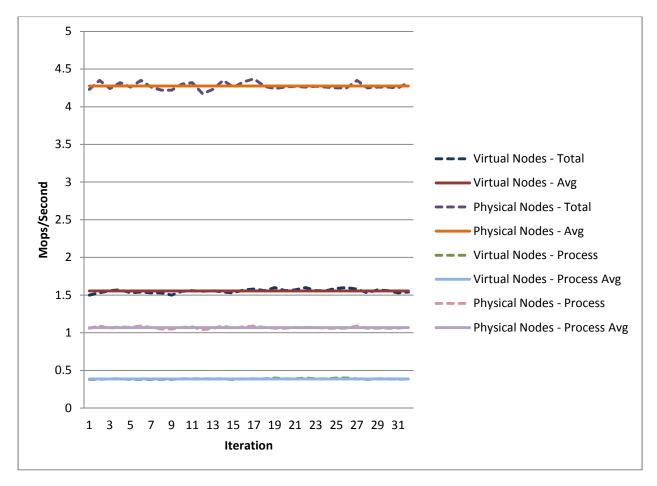


Figure A18. NAS Parallel Benchmark LU

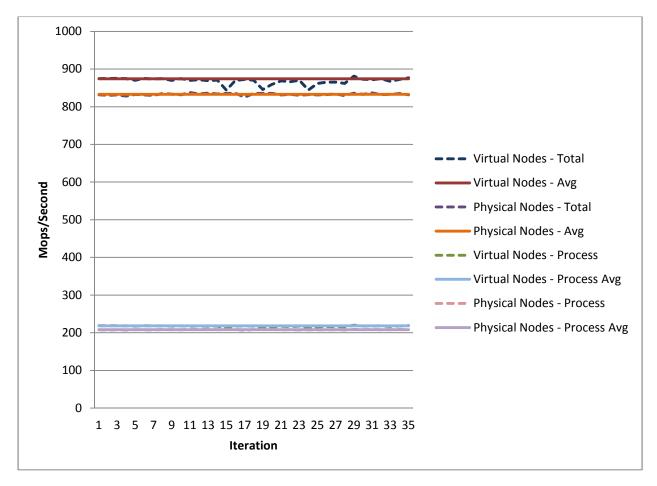


Figure A19. NAS Parallel Benchmarks MG

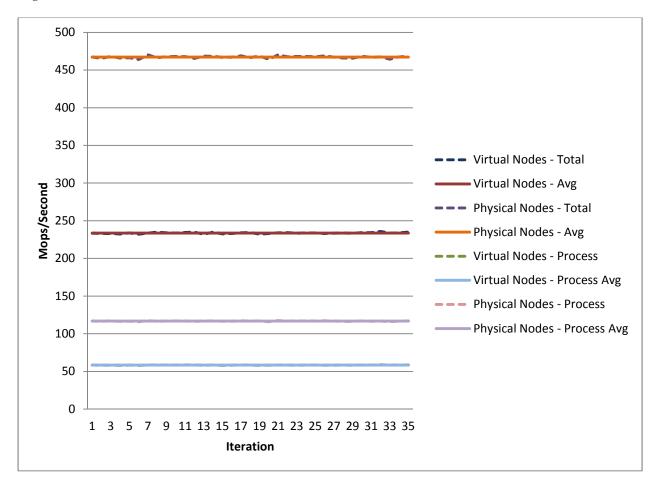


Figure A20. NAS Parallel Benchmarks SP

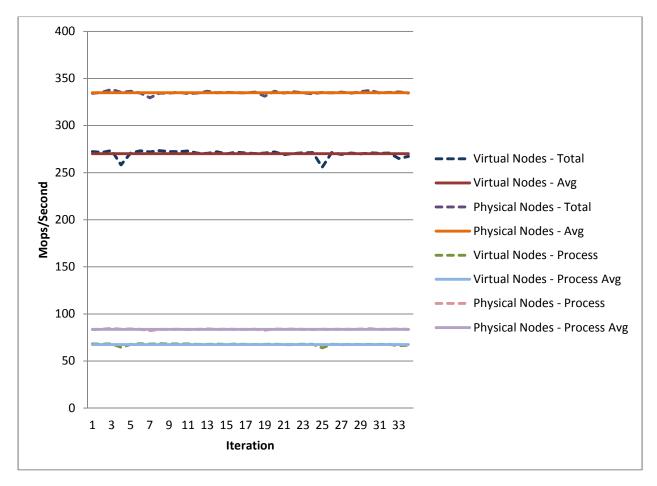


Figure A21. XHPL Benchmarks

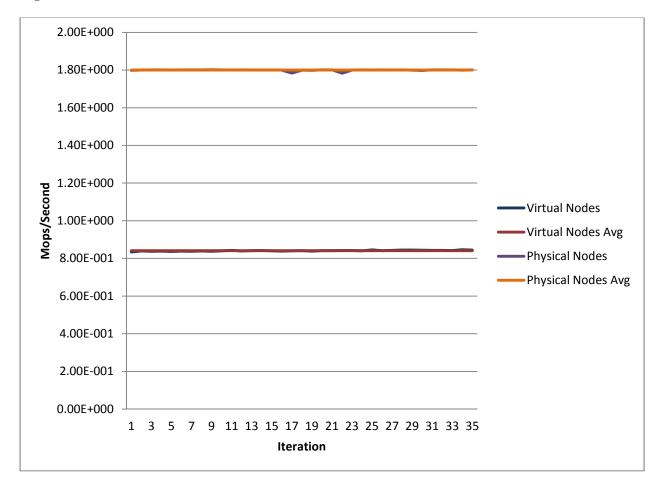


Figure A22. Filecp-lustre.pl

```
#!/usr/bin/perl
#filecp-lustre.pl
use File::Copy;
use Benchmark;
use Socket;
use Sys::Hostname;
open FILE, ">lustre1_4.txt" or die $!;
my $host = hostname();
print FILE $host,"\n";
$source = "/mnt/lustrefs/ubuntu-7.04-server-i386.iso";
$destination = "/tmp/ubuntu-7.04-server-i386.iso";
$starttime = new Benchmark;
$endtime = new Benchmark;
for ($count=0; $count <=35; $count++)
       t0 = new Benchmark;
       copy ($source, $destination) or die "File cannot be copied.";
       t1 = new Benchmark;
       td = timediff(t1, t0);
       ($time,$wallseconds)=split(\lambdas+/,timestr($td));
       print "$time","\n";
       print FILE $time,"\n";
       print split(\\s/,timestr(\$td),\$td),\\\n\\;
#
       unlink($destination);
#print "The code took:", $td, "\n";
$totaltime=timediff($endtime,$starttime);
close FILE;
```

#### Figure A23. Filecp-pvfs.pl

```
#!/usr/bin/perl
#filecp-pvfs.pl
use File::Copy;
use Benchmark;
use Socket;
use Sys::Hostname;
open FILE, ">pvfs1_4.txt" or die $!;
my $host = hostname();
print FILE $host,"\n";
$source = "/mnt/pvfs/ubuntu-7.04-server-i386.iso";
$destination = "/tmp/ubuntu-7.04-server-i386.iso";
$starttime = new Benchmark;
$endtime = new Benchmark;
for ($count=0; $count <= 35; $count++)
       $t0 = new Benchmark;
       copy ($source, $destination) or die "File cannot be copied.";
       $t1 = new Benchmark;
       td = timediff(t1, t0);
       (\text{stime},\text{swallseconds})=\text{split}(/\s+/,\text{timestr}(\text{std}));
       print "$time","\n";
       print FILE $time,"\n";
       print split(\\s/,timestr(\$td),\$td),\"\n";
#
       unlink($destination);
#print "The code took:", $td, "\n";
$totaltime=timediff($endtime,$starttime);
close FILE;
```

#### Figure A24. Filecp-nfs.pl

```
#!/usr/bin/perl
#filecp-nfs.pl
use File::Copy;
use Benchmark;
use Socket;
use Sys::Hostname;
open FILE, ">nfs1_4.txt" or die $!;
my $host = hostname();
print FILE $host,"\n";
$source = "/mnt/nfsmount/ubuntu-7.04-server-i386.iso";
$destination = "/tmp/ubuntu-7.04-server-i386.iso";
$starttime = new Benchmark;
$endtime = new Benchmark;
for ($count=0; $count <= 35; $count++)
       $t0 = new Benchmark;
       copy ($source, $destination) or die "File cannot be copied.";
       $t1 = new Benchmark;
       td = timediff(t1, t0);
       ($time,$wallseconds)=split(\lambdas+/,timestr($td));
       print "$time","\n";
       print FILE $time,"\n";
       print split(\\s/,timestr(\$td),\$td),\"\n";
#
       unlink($destination);
#print "The code took:", $td, "\n";
$totaltime=timediff($endtime,$starttime);
close FILE;
```

Figure A25. Filecp-Lustre\_sf.pl

```
#!/usr/bin/perl
#filecp-lustre_sf.pl
use File::Copy;
use Benchmark;
use Socket;
use Sys::Hostname;
open FILE, ">lustre_sf1_4.txt" or die $!;
my $host = hostname();
print FILE $host,"\n";
for ($outercount=0; $outercount<=34; $outercount++)
       $starttime = new Benchmark;
               for ($count=0; $count<=2047; $count++)
                       $sourcefilename="/tmp/smallfiles/"."file".".$count;
                       $destinationfilename="/mnt/lustrefs/lustre7/"."file".".$count;
                       print "$sourcefilename";
                       print " ";
                       print "$destinationfilename";
                       system("cp -f $sourcefilename $destinationfilename");
                       print "\n";
               }
$endtime= new Benchmark;
$td = timediff($endtime,$starttime);
(\text{time},\text{swallseconds})=\text{split}(/\s+/,\text{timestr}(\text{td}));
print "$time","\n";
print FILE $time,"\n";
close FILE;
```

#### Figure A26. Filecp-pvfs\_sf.pl

```
#!/usr/bin/perl
#filecp-pvfs_sf.pl
use File::Copy;
use Benchmark;
use Socket;
use Sys::Hostname;
open FILE, ">pvfs_sf1_4.txt" or die $!;
my $host = hostname();
print FILE $host,"\n";
for ($outercount=0; $outercount<=34; $outercount++)
       $starttime = new Benchmark;
               for ($count=0; $count<=2047; $count++)
                       $sourcefilename="/tmp/smallfiles/"."file".".$count;
                      $destinationfilename="/mnt/pvfs/lustre7/"."file".".$count;
                       print "$sourcefilename";
                       print " ";
                       print "$destinationfilename";
                       system("cp -f $sourcefilename $destinationfilename");
                       print "\n";
$endtime= new Benchmark;
$td = timediff($endtime,$starttime);
(\text{stime},\text{swallseconds})=\text{split}(/\s+/,\text{timestr}(\text{std}));
print "$time","\n";
print FILE $time,"\n";
close FILE;
```

Figure A27. Filecp-Lustre\_sf.pl

```
#!/usr/bin/perl
#filecp-nfs_sf.pl
use File::Copy;
use Benchmark;
use Socket;
use Sys::Hostname;
open FILE, ">nfs_sf1_4.txt" or die $!;
my $host = hostname();
print FILE $host,"\n";
for ($outercount=0; $outercount<=34; $outercount++)
       $starttime = new Benchmark;
               for ($count=0; $count<=2047; $count++)
                       $sourcefilename="/tmp/smallfiles/"."file".".$count;
                      $destinationfilename="/mnt/nfsmount/lustre7/"."file".".$count;
                       print "$sourcefilename";
                       print " ";
                       print "$destinationfilename";
                       system("cp -f $sourcefilename $destinationfilename");
                       print "\n";
               }
$endtime= new Benchmark;
$td = timediff($endtime,$starttime);
(\text{stime},\text{swallseconds})=\text{split}(/\s+/,\text{timestr}(\text{std}));
print "$time","\n";
print FILE $time,"\n";
close FILE;
```

Figure A28. Smallfiles-lustre.pl

```
#!/usr/bin/perl
#smallfiles.pl
use File::Copy;
use Benchmark;
use Socket;
use Sys::Hostname;
open FILE, ">sf-lustre1_4.txt" or die $!;
my $host = hostname();
print FILE $host,"\n";
#$delfiles = "/mnt/nfsmount/file*.txt";
for ($outercount=0; $outercount<=63;)</pre>
       $t0 = new Benchmark;
       for ($count=0; $count<=1023; $count++)
       system("dd if=/dev/urandom of=/mnt/lustrefs/file.$count bs=1024 count=1");
$t1 = new Benchmark;
td = timediff(t1,t0);
(\text{stime},\text{swallseconds})=\text{split}(/\s+/,\text{timestr}(\text{std}));
print "$time","\n";
print FILE $time,"\n";
select(undef,undef,undef,.250);
#unlink($delfiles);
$outercount++;
```

Figure A29. Smallfiles-pvfs.pl

```
#!/usr/bin/perl
#smallfiles.pl
use File::Copy;
use Benchmark;
use Socket;
use Sys::Hostname;
open FILE, ">sf-pvfs1_4.txt" or die $!;
my $host = hostname();
print FILE $host,"\n";
#$delfiles = "/mnt/nfsmount/file*.txt";
for ($outercount=0; $outercount<=63;)</pre>
       $t0 = new Benchmark;
       for ($count=0; $count<=1023; $count++)
       system("dd if=/dev/urandom of=/mnt/pvfs/file.$count bs=1024 count=1");
$t1 = new Benchmark;
td = timediff(t1,t0);
($time,$wallseconds)=split(\lambdas+/,timestr($td));
print "$time","\n";
print FILE $time,"\n";
select(undef,undef,undef,.250);
#unlink($delfiles);
$outercount++;
```

#### Figure A30. Smallfiles-nfs.pl

```
#!/usr/bin/perl
#smallfiles.pl
use File::Copy;
use Benchmark;
use Socket;
use Sys::Hostname;
open FILE, ">sf-nfs1_4.txt" or die $!;
my $host = hostname();
print FILE $host,"\n";
#$delfiles = "/mnt/nfsmount/file*.txt";
for ($outercount=0; $outercount<=63;)</pre>
       $t0 = new Benchmark;
       for ($count=0; $count<=1023; $count++)
       system("dd if=/dev/urandom of=/mnt/nfsmount/file.$count bs=1024 count=1");
$t1 = new Benchmark;
td = timediff(t1,t0);
(\text{stime},\text{swallseconds})=\text{split}(/\s+/,\text{timestr}(\text{std}));
print "$time","\n";
print FILE $time,"\n";
select(undef,undef,undef,.250);
#unlink($delfiles);
$outercount++;
```

Figure A31. HPLinpack benchmark input file

```
Innovative Computing Laboratory, University of Tennessee
HPL.out
           output file name (if any)
       device out (6=stdout,7=stderr,file)
6
       # of problems sizes (N)
1
4942 Ns
1
        # of NBs
16
     NBs
        PMAP process mapping (0=Row-,1=Column-major)
0
        # of process grids (P x Q)
1
1
      Ps
4
      Qs
16.0
         threshold
        # of panel fact
012
         PFACTs (0=left, 1=Crout, 2=Right)
        # of recursive stopping criterium
3
246
          NBMINs (>= 1)
        # of panels in recursion
2
        NDIVs
3
        # of recursive panel fact.
012
         RFACTs (0=left, 1=Crout, 2=Right)
1
        # of broadcast
0
        BCASTs (0=1rg,1=1rM,2=2rg,3=2rM,4=Lng,5=LnM)
        # of lookahead depth
1
0
        DEPTHs (>=0)
2
        SWAP (0=bin-exch,1=long,2=mix)
64
        swapping threshold
0
        L1 in (0=transposed,1=no-transposed) form
        U in (0=transposed,1=no-transposed) form
0
        Equilibration (0=no,1=yes)
1
        memory alignment in double (> 0)
8
```

Figure A32. Benchmarks.sh/Benchmarks-virt.sh

#!/bin/bash echo "cg" /home/glen/./submit-cg.pl

echo "mg" /home/glen/./submit-mg.pl

echo "is" /home/glen/./submit-is.pl

echo "ep" /home/glen/./submit-ep.pl

echo "lu" /home/glen/./submit-lu.pl

echo "ft" /home/glen/./submit-ft.pl

echo "sp" /home/glen/./submit-sp.pl

echo "bt" /home/glen/./submit-bt.pl

echo "xhpl" /home/glen/./submit-xhpl.pl

Figure A33. Script-Benchmark.sh

```
#!/bin/bash
#PBS -N bt.B.4
#nodes=4:ppn=1
#PBS -1
nodes = oscarnode 1. home.net + oscarnode 2. home.net + oscarnode 3. home.net + oscarnode 4. home.ne
#PBS -1 walltime=02:00:00
#PBS -r n
# -e stderr
# -o stdout
#PBS-V
echo $PBS_NODEFILE
cat $PBS_NODEFILE
#mpiexec /mnt/lustrefs/homes/glen/ep.A.2 -n 2
#/opt/mpich-1.2.7p1/bin/mpirun /mnt/lustrefs/homes/glen/ep.B.2 -v -machinefile
$PBS_NODEFILE -np 2
/opt/mpich-ch_p4-gcc-1.2.7/bin/mpirun -np 4 -machinefile $PBS_NODEFILE /home/glen/bt.B.4
```

```
Figure A34. Submit-Benchmark.pl
#!/usr/bin/perl
use File::Copy;
use Benchmark;
use Socket;
use Sys::Hostname;

for ($count=1;$count<=35;$count++)
{
    print $count;
    system("qsub script-cg");
    sleep (650);
}
```

#### Appendix B

#### File-System Benchmarks for physical cluster nodes

Figure B1. Lustre Large File Copy (1 of 1)
oscarnode1.home.net
45
44
44
45
44
45
45
44
44
45
45
44
45
45
44
45
45
44
44
44
44
44
45
45
44
44
44
45
45
44
45
45
44
44
4.4

*Figure B2*. Lustre Large File Copy (1 of 2) oscarnode1.home.net

*Figure B3*. Lustre Large File Copy (2 of 2) oscarnode2.home.net

```
Figure B4. Lustre Large File Copy (1 of 3) oscarnode1.home.net
69
69
69
69
69
70
68
69
68
70
69
69
68
68
70
71
68
69
70
70
70
69
69
69
70
69
70
69
70
69
69
```

```
Figure B5. Lustre Large File Copy (2 of 3) oscarnode2.home.net
69
69
69
69
69
69
69
70
69
69
69
69
68
69
70
71
69
69
69
70
70
68
69
68
71
69
71
70
70
70
69
```

```
Figure B6. Lustre Large File Copy (3 of 3) oscarnode3.home.net
69
70
69
69
69
69
69
69
69
69
69
69
70
68
70
71
68
69
69
70
71
69
69
70
70
70
70
70
70
69
70
```

```
Figure B7. Lustre Large File Copy (3 of 3) oscarnode3.home.net
69
70
69
69
69
69
69
69
69
69
69
69
70
68
70
71
68
69
69
70
71
69
69
70
70
70
70
70
70
```

```
Figure B8. Lustre Large File Copy (1 of 4) oscarnode1.home.net
89
89
89
90
90
89
90
90
91
90
92
90
91
90
89
90
89
92
91
89
91
91
91
90
91
89
91
90
89
89
90
```

```
Figure B9. Lustre Large File Copy (2 of 4) oscarnode2.home.net
90
89
89
91
90
89
89
90
90
91
91
91
91
90
90
90
89
91
90
88
90
90
90
90
90
89
91
89
90
89
90
```

```
Figure B10. Lustre Large File Copy (3 of 4) oscarnode3.home.net 89 90 89 91 90 89 90
```

```
Figure B11. Lustre Large File Copy (4 of 4) oscarnode4.home.net 90 89 90 90 90 89 90
```

# Figure B12. NFS Large File Copy (1 of 1) oscarnode1.home.net 47 47 48 47 47 47 47 47 47 47 47

#### Figure B13. NFS Large File Copy (1 of 2) oscarnode1.home.net

#### Figure B14. NFS Large File Copy (2 of 2) oscarnode2.home.net

```
Figure B15. NFS Large File Copy (1 of 3)
oscarnode1.home.net
69
68
68
68
68
72
68
69
66
68
68
68
68
69
68
68
69
69
67
68
68
69
69
68
67
70
71
69
```

```
Figure B16. NFS Large File Copy (2 of 3)
oscarnode2.home.net
69
69
66
68
67
72
68
67
70
68
67
69
68
69
68
68
67
68
68
67
69
```

```
Figure B17. NFS Large File Copy (3 of 3)
oscarnode3.home.net
69
69
66
68
67
71
67
68
69
68
67
68
67
69
69
68
69
68
68
68
69
69
70
67
68
69
```

## Figure B18. NFS Large File Copy (1 of 4) oscarnode1.home.net 91 86 86 86

#### Figure B19. NFS Large File Copy (2 of 4) oscarnode2.home.net

#### Figure B20. NFS Large File Copy (3 of 4) oscarnode3.home.net

### *Figure B21.* NFS Large File Copy (4 of 4) oscarnode4.home.net

```
Figure B22. PVFS Large File Copy (1 of 1) oscarnode1.home.net
51
49
50
49
49
49
49
49
49
50
50
49
49
50
49
49
50
50
49
49
50
50
49
49
49
50
49
49
50
50
49
50
```

Figure B23. PVFS Large File Copy (1 of 2) oscarnode1.home.net 

#### Figure B24. PFVS Large File Copy (2 of 2) oscarnode2.home.net

#### Figure B25. PVFS Large File Copy (1 of 3) oscarnode1.home.net

#### Figure B26. PVFS Large File Copy (2 of 3) oscarnode2.home.net

#### Figure B27. PVFS Large File Copy (3 of 3) oscarnode3.home.net

#### Figure B28. PVFS Large File Copy (1 of 4) oscarnode1.home.net

#### Figure B29. PVFS Large File Copy (2 of 4) oscarnode2.home.net

# Figure B30. PVFS Large File Copy (3 of 4) oscarnode3.home.net

# Figure B31. PVFS Large File Copy (4 of 4) oscarnode4.home.net 119

```
Figure B32. Lustre Small File Copy (1 of 1) oscarnode1.home.net
27
22
23
23
22
22
23
22
22
22
22
22
23
22
22
22
22
22
23
22
22
22
22
23
22
22
22
22
23
```

*Figure B33*. Lustre Small File Copy (1 of 2) oscarnode1.home.net

*Figure B34*. Lustre Small File Copy (2 of 2) oscarnode2.home.net

*Figure B35*. Lustre Small File Copy (1 of 3) oscarnode1.home.net 

## *Figure B36.* Lustre Small File Copy (2 of 3) oscarnode2.home.net

### *Figure B37*. Lustre Small File Copy (3 of 3) oscarnode3.home.net

### *Figure B38.* Lustre Small File Copy (1 of 4) oscarnode1.home.net

### *Figure B39.* Lustre Small File Copy (2 of 4) oscarnode2.home.net

### *Figure B40.* Lustre Small File Copy (3 of 4) oscarnode3.home.net

#### Figure B41. Lustre Small File Copy (4 of 4) oscarnode4.home.net

#### Figure B42. NFS Small File Copy (1 of 1) oscarnode1.home.net

#### Figure B43. NFS Small File Copy (1 of 2) oscarnode1.home.net

#### Figure B44. NFS Small File Copy (2 of 2) oscarnode2.home.net

#### Figure B45. NFS Small File Copy (1 of 3) oscarnode1.home.net

#### Figure B46. NFS Small File Copy (2 of 3) oscarnode2.home.net

#### Figure B47. NFS Small File Copy (3 of 3) oscarnode3.home.net

#### Figure B48. NFS Small File Copy (1 of 4) oscarnode1.home.net

#### Figure B49. NFS Small File Copy (2 of 4) oscarnode2.home.net

### *Figure B50.* NFS Small File Copy (3 of 4) oscarnode3.home.net

#### Figure B51. NFS Small File Copy (4 of 4) oscarnode4.home.net

### Figure B52. PVFS Small File Copy (1 of 1) oscarnode1.home.net

J4

#### Figure B53. PVFS Small File Copy (1 of 2) oscarnode1.home.net

#### Figure B54. PVFS Small File Copy (2 of 2) oscarnode2.home.net

#### Figure B55. PVFS Small File Copy (1 of 3) oscarnode1.home.net

#### Figure B56. PVFS Small File Copy (2 of 3) oscarnode2.home.net

### Figure B57. PVFS Small File Copy (3 of 3)

oscarnode3.home.net

0)

#### Figure B58. PVFS Small File Copy (1 of 4) oscarnode1.home.net

#### Figure B59. PVFS Small File Copy (2 of 4) oscarnode2.home.net

#### Figure B60. PVFS Small File Copy (3 of 4) oscarnode3.home.net

#### Figure B61. PVFS Small File Copy (4 of 4) oscarnode4.home.net

```
Figure B62. Lustre Small File Write (1 of 1)
oscarnode1.home.net
15
10
10
13
10
10
12
10
10
10
11
11
10
10
11
10
11
10
10
10
10
10
11
11
11
11
11
10
10
12
10
11
```

```
Figure B63. Lustre Small File Write (1 of 2)
oscarnode1.home.net
11
10
10
13
11
10
12
11
10
11
12
11
10
11
12
11
11
12
11
11
11
```

# Figure B64. Lustre Small File Write (2 of 2) oscarnode2.home.net 11 10 10 14 10 11 11 11

#### Figure B65. Lustre Small File Write (1 of 3) oscarnode1.home.net

#### Figure B66. Lustre Small File Write (2 of 3) oscarnode2.home.net

# *Figure B67.* Lustre Small File Write (3 of 3) oscarnode3.home.net

# *Figure B68.* Lustre Small File Write (1 of 4) oscarnode1.home.net

#### Figure B69. Lustre Small File Write (2 of 4) oscarnode2.home.net

# *Figure B70.* Lustre Small File Copy (3 of 4) oscarnode3.home.net

# Figure B71. Lustre Small File Write (4 of 4) oscarnode4.home.net 11 10 15 10 10 12 10 10

```
Figure B72. NFS Small File Write (1 of 1)
oscarnode1.home.net
12
10
10
10
11
11
10
11
11
11
11
10
10
11
11
10
10
10
11
13
11
10
10
10
10
10
```

#### Figure B73. NFS Small File Write (1 of 2) oscarnode1.home.net

#### Figure B74. NFS Small File Write (2 of 2) oscarnode2.home.net

### Figure B75. NFS Small File Write (1 of 3) oscarnode1.home.net

#### Figure B76. NFS Small File Write (2 of 3) oscarnode2.home.net

### Figure B77. NFS Small File Write (3 of 3) oscarnode3.home.net

# *Figure B78.* NFS Small File Write (1 of 4) oscarnode1.home.net

#### Figure B79. NFS Small File Write (2 of 4) oscarnode2.home.net

#### Figure B80. NFS Small File Write (3 of 4) oscarnode3.home.net

#### Figure B81. NFS Small File Write (4 of 4) oscarnode4.home.net

#### Figure B82. PVFS Small File Write (1 of 1) oscarnode1.home.net

#### Figure B83. PVFS Small File Write (1 of 2) oscarnode1.home.net

#### Figure B84. PVFS Small File Write (2 of 2) oscarnode2.home.net

#### Figure B85. PVFS Small File Write (1 of 3) oscarnode1.home.net

#### Figure B86. PVFS Small File Write (2 of 3) oscarnode2.home.net

# Figure B87. PVFS Small File Write (3 of 3) oscarnode3.home.net

# Figure B88. PVFS Small File Write (1 of 4) oscarnode1.home.net

# Figure B89. PVFS Small File Write (2 of 4) oscarnode2.home.net

#### Figure B90. PVFS Small File Write (3 of 4) oscarnode3.home.net

#### Figure B91. PVFS Small File Write (4 of 4) oscarnode4.home.net

#### **Appendix C**

#### File-System Benchmarks for virtual cluster nodes

```
Figure C2. Lustre Large File Copy (1 of 2) oscarnode5.home.net
73
70
71
71
71
68
69
70
68
70
70
70
76
73
71
72
69
79
68
69
72
71
71
67
70
72
67
72
71
73
68
68
```

Figure C3. Lustre Large File Copy (2 of 2) oscarnode6.home.net 

```
Figure C4. Lustre Large File Copy (1 of 3)
oscarnode5.home.net
118
109
108
101
103
99
94
119
117
99
93
92
97
98
101
96
92
92
95
99
99
98
100
96
104
100
94
99
```

# Figure C5. Lustre Large File Copy (2 of 3) oscarnode6.home.net 123 108 104

```
Figure C6. Lustre Large File Copy (3 of 3)
oscarnode7.home.net
87
89
89
97
95
102
104
94
82
92
92
101
107
103
99
99
100
104
99
100
106
99
99
100
100
108
94
104
97
94
94
```

```
Figure C7. Lustre Large File Copy (1 of 4)
oscarnode5.home.net
104
105
107
105
101
103
106
107
108
123
116
106
104
101
116
101
115
108
116
115
107
106
103
99
104
110
108
102
99
100
99
```

### Figure C8. Lustre Large File Copy (2 of 4) oscarnode6.home.net

```
Figure C9. Lustre Large File Copy (3 of 4)
oscarnode7.home.net
105
101
115
110
100
103
111
104
105
136
104
106
107
106
107
105
104
100
103
104
109
103
109
114
106
106
107
117
106
110
109
116
```

## Figure C10. Lustre Large File Copy (4 of 4) oscarnode8.home.net

```
Figure C11. NFS Large File Copy (1 of 1) oscarnode5.home.net
79
77
74
75
78
77
82
77
75
77
77
76
76
79
76
76
76
76
78
77
76
77
78
75
77
76
76
76
78
76
79
77
```

### Figure C12. NFS Large File Copy (1 of 2) oscarnode5.home.net

# Figure C13. NFS Large File Copy (2 of 2) oscarnode6.home.net

# Figure C14. NFS Large File Copy (1 of 3) oscarnode5.home.net 168 166 181 180 167 171

### Figure C15. NFS Large File Copy (2 of 3) oscarnode6.home.net

### Figure C16. NFS Large File Copy (3 of 3) oscarnode7.home.net

### Figure C17. NFS Large File Copy (1 of 4) oscarnode5.home.net

### Figure C18. NFS Large File Copy (2 of 4) oscarnode6.home.net

### Figure C19. NFS Large File Copy (3 of 4) oscarnode7.home.net

### Figure C20. NFS Large File Copy (4 of 4) oscarnode8.home.net

```
Figure C21. PVFS Large File Copy (1 of 1) oscarnode5.home.net
59
56
56
57
56
56
55
53
54
67
55
54
55
55
58
55
56
57
57
61
58
58
65
57
57
60
62
62
60
64
64
```

### Figure C22. PVFS Large File Copy (1 of 2) oscarnode5.home.net

### Figure C23. PVFS Large File Copy (1 of 3) oscarnode5.home.net

### Figure C24. PVFS Large File Copy (2 of 3) oscarnode6.home.net

### Figure C25. PVFS Large File Copy (3 of 3) oscarnode7.home.net

### Figure C26. PVFS Large File Copy (1 of 4) oscarnode5.home.net

```
Figure C27. PVFS Large File Copy (2 of 4)
oscarnode6.home.net
102
101
103
111
111
107
111
106
108
108
105
104
104
118
114
105
106
111
114
110
107
102
99
110
107
114
109
109
110
110
108
101
```

### Figure C28. PVFS Large File Copy (3 of 4) oscarnode7.home.net

### Figure C29. PVFS Large File Copy (4 of 4) oscarnode8.home.net

### Figure C30. Lustre Small File Copy (1 of 1) oscarnode5.home.net

Figure C31. Lustre Small File Copy (1 of 2) oscarnode5.home.net 28 29 26

### Figure C34. Lustre Small File Copy (2 of 2) oscarnode6.home.net

*Figure C35*. Lustre Small File Copy (1 of 3) oscarnode5.home.net

```
Figure C36. Lustre Small File Copy (2 of 3)
oscarnode6.home.net
45
33
48
48
48
48
47
48
46
44
46
47
45
47
46
48
48
48
47
48
48
48
48
48
48
47
49
47
```

### Figure C37. Lustre Small File Copy (3 of 3) oscarnode7.home.net

```
Figure C38. Lustre Small File Copy (1 of 4)
oscarnode5.home.net
49
123
119
108
98
98
97
97
93
82
77
74
73
68
68
66
66
65
65
64
65
65
65
65
64
65
65
```

```
Figure C39. Lustre Small File Copy (2 of 4) oscarnode6.home.net
50
146
120
106
99
97
97
97
90
84
77
74
71
70
69
68
67
65
65
67
66
66
65
66
66
65
65
66
67
66
```

```
Figure C40. Lustre Small File Copy (3 of 4) oscarnode7.home.net
75
120
120
107
96
94
96
97
89
83
78
73
73
69
68
68
65
66
66
65
65
65
67
65
65
65
66
66
65
64
66
```

```
Figure C41. Lustre Small File Copy (4 of 4)
oscarnode8.home.net
49
126
120
109
98
99
99
97
92
82
80
75
72
70
70
68
67
65
66
67
65
66
66
65
67
66
66
65
66
65
65
```

### Figure C42. NFS Small File Copy (1 of 1) oscarnode5.home.net

# Figure C43. NFS Small File Copy (1 of 2) oscarnode5.home.net

#### Figure C44. NFS Small File Copy (2 of 2) oscarnode6.home.net

#### Figure C45. NFS Small File Copy (1 of 3) oscarnode5.home.net

#### Figure C46. NFS Small File Copy (2 of 3) oscarnode6.home.net

#### Figure C47. NFS Small File Copy (3 of 3) oscarnode7.home.net

#### Figure C48. NFS Small File Copy (1 of 4) oscarnode5.home.net

# Figure C49. NFS Small File Copy (2 of 4) oscarnode6.home.net

#### Figure C50. NFS Small File Copy (3 of 4) oscarnode7.home.net

#### Figure C51. NFS Small File Copy (4 of 4) oscarnode8.home.net

```
Figure C52. Lustre Small File Write (1 of 1)
oscarnode5.home.net
15
10
10
14
10
10
12
12
10
11
12
10
11
11
12
11
11
12
12
11
11
12
11
11
11
12
11
11
12
11
11
```

Figure C53. Luster Small File Write (1 of 2) oscarnode5.home.net

# *Figure C54*. Lustre Small File Write (2 of 2) oscarnode6.home.net

#### Figure C55. Lustre Small File Write (1 of 3) oscarnode5.home.net

#### Figure C56. Lustre Small File Write (2 of 3) oscarnode6.home.net

# *Figure C57.* Lustre Small File Write (3 of 3) oscarnode7.home.net

#### Figure C58. Lustre Small File Write (1 of 4) oscarnode5.home.net

#### Figure C59. Lustre Small File Write (2 of 4) oscarnode6.home.net

#### Figure C60. Lustre Small File Write (3 of 4) oscarnode7.home.net

#### Figure C61. Lustre Small File Write (4 of 4) oscarnode8.home.net

# Figure C62. NFS Small File Write (1 of 1) oscarnode5.home.net 15

# *Figure C63*. NFS Small File Write (1 of 2) oscarnode5.home.net

#### Figure C64. NFS Small File Write (2 of 2) oscarnode6.home.net

# Figure C65. NFS Small File Write (1 of 3) oscarnode5.home.net

#### Figure C66. NFS Small File Write (2 of 3) oscarnode6.home.net

#### Figure C67. NFS Small File Write (3 of 3) oscarnode7.home.net

#### Figure C68. NFS Small File Write (1 of 4) oscarnode5.home.net

#### Figure C69. NFS Small File Write (2 of 4) oscarnode6.home.net

# Figure C70. NFS Small File Write (3 of 4) oscarnode7.home.net 28 28 28

#### Figure C71. NFS Small File Write (4 of 4) oscarnode8.home.net

#### Figure C72. PVFS Small File Write (1 of 1) oscarnode5.home.net

#### Figure C73. PVFS Small File Write (1 of 2) oscarnode5.home.net

#### Figure C74. PVFS Small File Write (2 of 2) oscarnode6.home.net

#### Figure C75. PVFS Small File Write (1 of 3) oscarnode5.home.net

#### Figure C76. PVFS Small File Write (2 of 3) oscarnode6.home.net

#### Figure C77. PVFS Small File Write (3 of 3) oscarnode7.home.net

# Figure C78. PVFS Small File Write (1 of 4) oscarnode5.home.net 60 63 64 65 65 70

```
Figure C79. PVFS Small File Write (2 of 4)
oscarnode6.home.net
60
65
64
67
66
61
64
69
67
69
66
70
73
71
70
72
70
70
71
70
70
72
68
69
68
69
70
69
67
63
69
```

#### Figure C80. PVFS Small File Copy (3 of 4) oscarnode7.home.net

#### Figure C81. PVFS Small File Write (4 of 4) oscarnode8.home.net

#### Appendix D

#### **HPC** Benchmarks

Figure D1. NAS Parallel Benchmark BT

Virtual Nodes - Total	Virtual Nodes - Avg	Physical Nodes - Total	Physical Nodes - Avg	Virtual Nodes - Process	Virtual Nodes - Process Avg	Physical Nodes - Process	Physical Nodes - Process Avg
698.56	694.6117143	630.85	629.9182857	174.64	173.6531	157.71	157.479
701.82	694.6117143	624.77	629.9182857	175.45	173.6531	156.19	157.479
697.93	694.6117143	625.27	629.9182857	174.48	173.6531	156.32	157.479
697.86	694.6117143	626.73	629.9182857	174.47	173.6531	156.68	157.479
700.54	694.6117143	629.95	629.9182857	175.13	173.6531	157.49	157.479
700.5	694.6117143	631.66	629.9182857	175.12	173.6531	157.91	157.479
700.6	694.6117143	629.81	629.9182857	175.15	173.6531	157.45	157.479
699.33	694.6117143	624.53	629.9182857	174.83	173.6531	156.13	157.479
700.12	694.6117143	633.9	629.9182857	175.03	173.6531	158.48	157.479
701.49	694.6117143	627.81	629.9182857	175.37	173.6531	156.95	157.479
688.88	694.6117143	626.22	629.9182857	172.22	173.6531	156.55	157.479
689.17	694.6117143	630.43	629.9182857	172.29	173.6531	157.61	157.479
687.71	694.6117143	631.45	629.9182857	171.93	173.6531	157.86	157.479
688.54	694.6117143	626.52	629.9182857	172.14	173.6531	156.63	157.479
689.1	694.6117143	629.84	629.9182857	172.28	173.6531	157.46	157.479
691.47	694.6117143	629	629.9182857	172.87	173.6531	157.25	157.479
689.99	694.6117143	630.31	629.9182857	172.5	173.6531	157.58	157.479
689.09	694.6117143	627.94	629.9182857	172.27	173.6531	156.99	157.479
694.33	694.6117143	633.85	629.9182857	173.58	173.6531	158.46	157.479
688.02	694.6117143	630.59	629.9182857	172.01	173.6531	157.65	157.479
690.82	694.6117143	630.71	629.9182857	172.71	173.6531	157.68	157.479
691.85	694.6117143	628.58	629.9182857	172.96	173.6531	157.15	157.479
693.08	694.6117143	631.56	629.9182857	173.27	173.6531	157.89	157.479
692.55	694.6117143	631.24	629.9182857	173.14	173.6531	157.81	157.479
689.45	694.6117143	632.67	629.9182857	172.36	173.6531	158.17	157.479
694.54	694.6117143	628.89	629.9182857	173.64	173.6531	157.22	157.479
690.96	694.6117143	627.32	629.9182857	172.74	173.6531	156.83	157.479
694.29	694.6117143	631.84	629.9182857	173.57	173.6531	157.96	157.479
694.38	694.6117143	633.07	629.9182857	173.6	173.6531	158.27	157.479
693.04	694.6117143	630.38	629.9182857	173.26	173.6531	157.59	157.479
698.45	694.6117143	632.05	629.9182857	174.61	173.6531	158.01	157.479
696.78	694.6117143	630.92	629.9182857	174.2	173.6531	157.73	157.479
697.51	694.6117143	633.68	629.9182857	174.38	173.6531	158.42	157.479
700.09	694.6117143	631.3	629.9182857	175.02	173.6531	157.82	157.479
698.57	694.6117143	631.5	629.9182857	174.64	173.6531	157.88	157.479

Figure D2. NAS Parallel Benchmarks CG

Virtual Nodes - Total	Virtual Nodes - Avg	Physical Nodes - Total	Physical Nodes - Avg	Virtual Nodes - Process	Virtual Nodes - Process Avg	Physical Nodes - Process	Physical Nodes - Process Avg
67.49	67.70176471	92.61	91.76857143	16.87	16.925	23.15	22.9417
67.43	67.70176471	92.2	91.76857143	16.86	16.925	23.05	22.9417
67.45	67.70176471	92.36	91.76857143	16.86	16.925	23.09	22.9417
67.66	67.70176471	92.53	91.76857143	16.91	16.925	23.13	22.9417
67.69	67.70176471	91.16	91.76857143	16.92	16.925	22.79	22.9417
67.68	67.70176471	92.08	91.76857143	16.92	16.925	23.02	22.9417
67.72	67.70176471	92.68	91.76857143	16.93	16.925	23.17	22.9417
67.81	67.70176471	92.36	91.76857143	16.95	16.925	23.09	22.9417
67.91	67.70176471	92.71	91.76857143	16.98	16.925	23.18	22.9417
67.94	67.70176471	91.76	91.76857143	16.99	16.925	22.94	22.9417
67.68	67.70176471	90.05	91.76857143	16.92	16.925	22.51	22.9417
67.65	67.70176471	89.66	91.76857143	16.91	16.925	22.42	22.9417
67.73	67.70176471	91.96	91.76857143	16.93	16.925	22.99	22.9417
67.7	67.70176471	91.58	91.76857143	16.93	16.925	22.89	22.9417
67.68	67.70176471	92.07	91.76857143	16.92	16.925	23.02	22.9417
67.73	67.70176471	92.14	91.76857143	16.93	16.925	23.03	22.9417
67.66	67.70176471	92.02	91.76857143	16.91	16.925	23	22.9417
67.72	67.70176471	92.07	91.76857143	16.93	16.925	23.02	22.9417
67.77	67.70176471	89.54	91.76857143	16.94	16.925	22.39	22.9417
67.75	67.70176471	91.62	91.76857143	16.94	16.925	22.9	22.9417
67.83	67.70176471	91.23	91.76857143	16.96	16.925	22.81	22.9417
67.82	67.70176471	92.69	91.76857143	16.96	16.925	23.17	22.9417
67.7	67.70176471	92.31	91.76857143	16.92	16.925	23.08	22.9417
67.76	67.70176471	92.38	91.76857143	16.94	16.925	23.09	22.9417
67.8	67.70176471	92.44	91.76857143	16.95	16.925	23.11	22.9417
67.64	67.70176471	91.25	91.76857143	16.91	16.925	22.81	22.9417
67.76	67.70176471	89.67	91.76857143	16.94	16.925	22.42	22.9417
67.23	67.70176471	92.59	91.76857143	16.81	16.925	23.15	22.9417
67.74	67.70176471	91.87	91.76857143	16.93	16.925	22.97	22.9417
67.73	67.70176471	92.15	91.76857143	16.93	16.925	23.04	22.9417
67.88	67.70176471	89.39	91.76857143	16.97	16.925	22.35	22.9417
67.74	67.70176471	92.02	91.76857143	16.94	16.925	23	22.9417
67.66	67.70176471	92.16	91.76857143	16.91	16.925	23.04	22.9417
67.72	67.70176471	92.13	91.76857143	16.93	16.925	23.03	22.9417

Figure D3. NAS Parallel Benchmarks EP

Virtual Nodes - Total	Virtual Nodes - Avg	Physical Nodes - Total	Physical Nodes - Avg	Virtual Nodes - Process	Virtual Nodes - Process Avg	Physical Nodes - Process	Nodes - Process Avg
41.66	42.085	25.34	25.32914286	10.42	10.52058824	6.34	6.33343
42.23	42.085	25.36	25.32914286	10.56	10.52058824	6.34	6.33343
42.13	42.085	25.36	25.32914286	10.53	10.52058824	6.34	6.33343
42.13	42.085	25.4	25.32914286	10.53	10.52058824	6.35	6.33343
42.06	42.085	25.44	25.32914286	10.52	10.52058824	6.36	6.33343
42.21	42.085	25.35	25.32914286	10.55	10.52058824	6.34	6.33343
42.08	42.085	25.35	25.32914286	10.52	10.52058824	6.34	6.33343
42.18	42.085	25.36	25.32914286	10.54	10.52058824	6.34	6.33343
42.15	42.085	25.36	25.32914286	10.54	10.52058824	6.34	6.33343
42.09	42.085	25.42	25.32914286	10.52	10.52058824	6.36	6.33343
42	42.085	25.44	25.32914286	10.5	10.52058824	6.36	6.33343
42.13	42.085	25.36	25.32914286	10.53	10.52058824	6.34	6.33343
42.12	42.085	25.4	25.32914286	10.53	10.52058824	6.35	6.33343
42.03	42.085	25.35	25.32914286	10.51	10.52058824	6.34	6.33343
41.9	42.085	25.34	25.32914286	10.48	10.52058824	6.33	6.33343
42.24	42.085	25.44	25.32914286	10.56	10.52058824	6.36	6.33343
41.98	42.085	25.44	25.32914286	10.49	10.52058824	6.36	6.33343
42.08	42.085	25.4	25.32914286	10.52	10.52058824	6.35	6.33343
42.06	42.085	25.35	25.32914286	10.51	10.52058824	6.34	6.33343
42.14	42.085	23.74	25.32914286	10.53	10.52058824	5.94	6.33343
42.06	42.085	25.35	25.32914286	10.51	10.52058824	6.34	6.33343
41.9	42.085	25.42	25.32914286	10.48	10.52058824	6.36	6.33343
42.15	42.085	25.34	25.32914286	10.54	10.52058824	6.34	6.33343
42.01	42.085	25.34	25.32914286	10.5	10.52058824	6.33	6.33343
42.29	42.085	25.36	25.32914286	10.57	10.52058824	6.34	6.33343
42.02	42.085	25.36	25.32914286	10.5	10.52058824	6.34	6.33343
42.05	42.085	25.35	25.32914286	10.51	10.52058824	6.34	6.33343
41.98	42.085	25.35	25.32914286	10.5	10.52058824	6.34	6.33343
42.12	42.085	25.42	25.32914286	10.53	10.52058824	6.35	6.33343
42.33	42.085	25.35	25.32914286	10.58	10.52058824	6.34	6.33343
42.03	42.085	25.36	25.32914286	10.51	10.52058824	6.34	6.33343
42.24	42.085	25.44	25.32914286	10.56	10.52058824	6.36	6.33343
41.98	42.085	25.34	25.32914286	10.49	10.52058824	6.34	6.33343
42.13	42.085	25.35	25.32914286	10.53	10.52058824	6.34	6.33343
42.07	42.085	25.39	25.32914286	10.52	10.52058824	6.35	6.33343

Figure D4. NAS Parallel Benchmarks FT

Virtual Nodes - Total	Virtual Nodes - Avg	Physical Nodes - Total	Physical Nodes - Avg	Virtual Nodes - Process	Virtual Nodes - Process Avg	Physical Nodes - Process	Physical Nodes - Process Avg
55.79	55.10558824	117.98	116.3662857	13.95	13.77647	29.49	29.0911
55.98	55.10558824	116.14	116.3662857	14	13.77647	29.04	29.0911
56.28	55.10558824	112.3	116.3662857	14.07	13.77647	28.08	29.0911
56.12	55.10558824	115.63	116.3662857	14.03	13.77647	28.91	29.0911
56.09	55.10558824	99.04	116.3662857	14.02	13.77647	24.76	29.0911
54.3	55.10558824	115.87	116.3662857	13.57	13.77647	28.97	29.0911
54.36	55.10558824	114.01	116.3662857	13.59	13.77647	28.5	29.0911
54.45	55.10558824	103.51	116.3662857	13.61	13.77647	25.88	29.0911
55.42	55.10558824	110.11	116.3662857	13.86	13.77647	27.53	29.0911
54.37	55.10558824	113.95	116.3662857	13.59	13.77647	28.49	29.0911
55.34	55.10558824	122.38	116.3662857	13.84	13.77647	30.59	29.0911
54.43	55.10558824	119.5	116.3662857	13.61	13.77647	29.87	29.0911
54.3	55.10558824	105.02	116.3662857	13.58	13.77647	26.25	29.0911
54.28	55.10558824	115.22	116.3662857	13.57	13.77647	28.8	29.0911
54.25	55.10558824	115.12	116.3662857	13.56	13.77647	28.78	29.0911
54.38	55.10558824	109.35	116.3662857	13.6	13.77647	27.34	29.0911
55.47	55.10558824	127.56	116.3662857	13.87	13.77647	31.89	29.0911
55.32	55.10558824	128.85	116.3662857	13.83	13.77647	32.21	29.0911
55.25	55.10558824	123	116.3662857	13.81	13.77647	30.75	29.0911
55.2	55.10558824	122.85	116.3662857	13.8	13.77647	30.71	29.0911
55.22	55.10558824	109.26	116.3662857	13.8	13.77647	27.31	29.0911
55.11	55.10558824	115.27	116.3662857	13.78	13.77647	28.82	29.0911
55.19	55.10558824	125.71	116.3662857	13.8	13.77647	31.43	29.0911
56.2	55.10558824	114.19	116.3662857	14.05	13.77647	28.55	29.0911
55.84	55.10558824	116.85	116.3662857	13.96	13.77647	29.21	29.0911
56.15	55.10558824	118.6	116.3662857	14.04	13.77647	29.65	29.0911
55.94	55.10558824	116.96	116.3662857	13.99	13.77647	29.24	29.0911
55.81	55.10558824	125	116.3662857	13.95	13.77647	31.25	29.0911
56.01	55.10558824	118.39	116.3662857	14	13.77647	29.6	29.0911
56.06	55.10558824	117.2	116.3662857	14.01	13.77647	29.3	29.0911
56.17	55.10558824	117.46	116.3662857	14.04	13.77647	29.36	29.0911
56.36	55.10558824	108.52	116.3662857	14.09	13.77647	27.13	29.0911
54.66	55.10558824	120.49	116.3662857	13.66	13.77647	30.12	29.0911
54.57	55.10558824	125.26	116.3662857	13.64	13.77647	31.31	29.0911
54.66	55.10558824	116.27	116.3662857	13.66	13.77647	29.07	29.0911

Figure D5. NAS Parallel Benchmarks IS

Virtual Nodes - Total	Virtual Nodes - Avg	Physical Nodes - Total	Physical Nodes - Avg	Virtual Nodes - Process	Virtual Nodes - Process Avg	Physical Nodes - Process	Physical Nodes - Process Avg
1.5	1.554375	4.23	4.274285714	0.38	0.388125	1.06	1.068
1.53	1.554375	4.35	4.274285714	0.38	0.388125	1.09	1.068
1.56	1.554375	4.24	4.274285714	0.39	0.388125	1.06	1.068
1.57	1.554375	4.32	4.274285714	0.39	0.388125	1.08	1.068
1.53	1.554375	4.26	4.274285714	0.38	0.388125	1.07	1.068
1.54	1.554375	4.35	4.274285714	0.38	0.388125	1.09	1.068
1.53	1.554375	4.26	4.274285714	0.38	0.388125	1.07	1.068
1.53	1.554375	4.22	4.274285714	0.38	0.388125	1.05	1.068
1.5	1.554375	4.22	4.274285714	0.38	0.388125	1.05	1.068
1.55	1.554375	4.3	4.274285714	0.39	0.388125	1.07	1.068
1.56	1.554375	4.32	4.274285714	0.39	0.388125	1.08	1.068
1.55	1.554375	4.17	4.274285714	0.39	0.388125	1.04	1.068
1.56	1.554375	4.23	4.274285714	0.39	0.388125	1.06	1.068
1.54	1.554375	4.35	4.274285714	0.39	0.388125	1.09	1.068
1.53	1.554375	4.26	4.274285714	0.38	0.388125	1.06	1.068
1.57	1.554375	4.33	4.274285714	0.39	0.388125	1.08	1.068
1.58	1.554375	4.37	4.274285714	0.39	0.388125	1.09	1.068
1.56	1.554375	4.27	4.274285714	0.39	0.388125	1.07	1.068
1.6	1.554375	4.24	4.274285714	0.4	0.388125	1.06	1.068
1.56	1.554375	4.26	4.274285714	0.39	0.388125	1.06	1.068
1.57	1.554375	4.27	4.274285714	0.39	0.388125	1.07	1.068
1.6	1.554375	4.26	4.274285714	0.4	0.388125	1.07	1.068
1.56	1.554375	4.27	4.274285714	0.39	0.388125	1.07	1.068
1.56	1.554375	4.26	4.274285714	0.39	0.388125	1.06	1.068
1.59	1.554375	4.25	4.274285714	0.4	0.388125	1.06	1.068
1.6	1.554375	4.25	4.274285714	0.4	0.388125	1.06	1.068
1.58	1.554375	4.35	4.274285714	0.39	0.388125	1.09	1.068
1.53	1.554375	4.25	4.274285714	0.38	0.388125	1.06	1.068
1.57	1.554375	4.26	4.274285714	0.39	0.388125	1.06	1.068
1.56	1.554375	4.26	4.274285714	0.39	0.388125	1.06	1.068
1.53	1.554375	4.25	4.274285714	0.38	0.388125	1.06	1.068
1.54	1.554375	4.32	4.274285714	0.39	0.388125	1.08	1.068

Figure D6. NAS Parallel Benchmarks LU

Virtual Nodes - Total	Virtual Nodes - Avg	Physical Nodes - Total	Physical Nodes - Avg	Virtual Nodes - Process	Virtual Nodes - Process Avg	Physical Nodes - Process	Physical Nodes - Process Avg
874.66	873.986	831.76	832.7982857	218.67	218.4963	207.94	208.199
874.82	873.986	830	832.7982857	218.71	218.4963	207.5	208.199
875.48	873.986	831.14	832.7982857	218.87	218.4963	207.78	208.199
874.34	873.986	828.14	832.7982857	218.58	218.4963	207.03	208.199
870.11	873.986	832.39	832.7982857	217.53	218.4963	208.1	208.199
875.02	873.986	831.08	832.7982857	218.76	218.4963	207.77	208.199
873.7	873.986	829.79	832.7982857	218.43	218.4963	207.45	208.199
874.59	873.986	834.92	832.7982857	218.65	218.4963	208.73	208.199
869.95	873.986	833.44	832.7982857	217.49	218.4963	208.36	208.199
874.36	873.986	831.49	832.7982857	218.59	218.4963	207.87	208.199
870.02	873.986	837.55	832.7982857	217.51	218.4963	209.39	208.199
871.45	873.986	833.59	832.7982857	217.86	218.4963	208.4	208.199
869.14	873.986	836.5	832.7982857	217.28	218.4963	209.13	208.199
870.91	873.986	833.86	832.7982857	217.73	218.4963	208.46	208.199
844.06	873.986	835.38	832.7982857	211.01	218.4963	208.85	208.199
870	873.986	835.74	832.7982857	217.5	218.4963	208.93	208.199
872.51	873.986	825.16	832.7982857	218.13	218.4963	206.29	208.199
870.08	873.986	835.44	832.7982857	217.52	218.4963	208.86	208.199
845.64	873.986	835.62	832.7982857	211.41	218.4963	208.9	208.199
859.44	873.986	835.54	832.7982857	214.86	218.4963	208.89	208.199
868.62	873.986	831.13	832.7982857	217.16	218.4963	207.78	208.199
866.1	873.986	832.8	832.7982857	216.53	218.4963	208.2	208.199
870.45	873.986	830.46	832.7982857	217.61	218.4963	207.61	208.199
844.47	873.986	832.06	832.7982857	211.12	218.4963	208.01	208.199
861.68	873.986	831.18	832.7982857	215.42	218.4963	207.8	208.199
865.13	873.986	831.87	832.7982857	216.28	218.4963	207.97	208.199
865.18	873.986	833.76	832.7982857	216.29	218.4963	208.44	208.199
861.77	873.986	829.58	832.7982857	215.44	218.4963	207.39	208.199
880.97	873.986	835.86	832.7982857	220.24	218.4963	208.97	208.199
872.96	873.986	832.97	832.7982857	218.24	218.4963	208.24	208.199
871.58	873.986	836.53	832.7982857	217.9	218.4963	209.13	208.199
874.41	873.986	831.87	832.7982857	218.6	218.4963	207.97	208.199
867.29	873.986	832.77	832.7982857	216.82	218.4963	208.19	208.199
872.42	873.986	834.91	832.7982857	218.1	218.4963	208.73	208.199
876.46	873.986	831.66	832.7982857	219.11	218.4963	207.91	208.199

Figure D7. NAS Parallel Benchmarks MG

Virtual Nodes - Total	Virtual Nodes - Avg	Physical Nodes - Total	Physical Nodes - Avg	Virtual Nodes - Process	Virtual Nodes - Process Avg	Physical Nodes - Process	Physical Nodes - Process Avg
233.53	233.6967647	467.19	467.2228571	58.38	58.4247	116.8	116.806
232.95	233.6967647	465.47	467.2228571	58.24	58.4247	116.37	116.806
232.98	233.6967647	467.97	467.2228571	58.25	58.4247	116.99	116.806
232.42	233.6967647	465.89	467.2228571	58.11	58.4247	116.47	116.806
233.81	233.6967647	466.36	467.2228571	58.45	58.4247	116.59	116.806
232.26	233.6967647	464.05	467.2228571	58.07	58.4247	116.01	116.806
233.72	233.6967647	470.19	467.2228571	58.43	58.4247	117.55	116.806
234.86	233.6967647	466.08	467.2228571	58.72	58.4247	116.52	116.806
234.17	233.6967647	467.81	467.2228571	58.54	58.4247	116.95	116.806
233.67	233.6967647	468.02	467.2228571	58.42	58.4247	117.01	116.806
234.4	233.6967647	467.76	467.2228571	58.6	58.4247	116.94	116.806
234.96	233.6967647	465.56	467.2228571	58.74	58.4247	116.39	116.806
231.95	233.6967647	468.54	467.2228571	57.99	58.4247	117.14	116.806
234.85	233.6967647	468.24	467.2228571	58.71	58.4247	117.06	116.806
232.56	233.6967647	466.75	467.2228571	58.14	58.4247	116.69	116.806
233.01	233.6967647	467	467.2228571	58.25	58.4247	116.75	116.806
234.11	233.6967647	468.84	467.2228571	58.53	58.4247	117.21	116.806
234.18	233.6967647	466.57	467.2228571	58.54	58.4247	116.64	116.806
232.12	233.6967647	468.38	467.2228571	58.03	58.4247	117.09	116.806
233.07	233.6967647	464.33	467.2228571	58.27	58.4247	116.08	116.806
234.07	233.6967647	470.27	467.2228571	58.52	58.4247	117.57	116.806
234.29	233.6967647	467.55	467.2228571	58.57	58.4247	116.89	116.806
233.39	233.6967647	468.01	467.2228571	58.35	58.4247	117	116.806
233.81	233.6967647	467.97	467.2228571	58.45	58.4247	116.99	116.806
233.8	233.6967647	467.71	467.2228571	58.45	58.4247	116.93	116.806
232.93	233.6967647	468.92	467.2228571	58.23	58.4247	117.23	116.806
233.49	233.6967647	467.08	467.2228571	58.37	58.4247	116.77	116.806
233.62	233.6967647	465.92	467.2228571	58.4	58.4247	116.48	116.806
233.47	233.6967647	465.72	467.2228571	58.37	58.4247	116.43	116.806
234.24	233.6967647	468.34	467.2228571	58.56	58.4247	117.09	116.806
234.26	233.6967647	466.78	467.2228571	58.57	58.4247	116.7	116.806
235.69	233.6967647	467.33	467.2228571	58.92	58.4247	116.83	116.806
233.55	233.6967647	464.47	467.2228571	58.39	58.4247	116.12	116.806
234.16	233.6967647	467.35	467.2228571	58.54	58.4247	116.84	116.806
234.87	233.6967647	468.38	467.2228571	58.72	58.4247	117.09	116.806

Figure D8. NAS Parallel Benchmarks SP

Virtual Nodes - Total	Virtual Nodes - Avg	Physical Nodes - Total	Physical Nodes - Avg	Virtual Nodes - Process	Virtual Nodes - Process Avg	Physical Nodes - Process	Physical Nodes - Process Avg
272.43	270.1773529	334.15	334.8657143	68.11	67.54470588	83.54	83.72617647
271.43	270.1773529	335.36	334.8657143	67.86	67.54470588	83.84	83.72617647
273.52	270.1773529	338.31	334.8657143	68.38	67.54470588	84.58	83.72617647
258.23	270.1773529	335.38	334.8657143	64.56	67.54470588	83.84	83.72617647
270.88	270.1773529	336.37	334.8657143	67.72	67.54470588	84.09	83.72617647
273.2	270.1773529	334.5	334.8657143	68.3	67.54470588	83.62	83.72617647
271.97	270.1773529	329.38	334.8657143	67.99	67.54470588	82.35	83.72617647
273.35	270.1773529	334.39	334.8657143	68.34	67.54470588	83.6	83.72617647
272.44	270.1773529	334.36	334.8657143	68.11	67.54470588	83.59	83.72617647
272.43	270.1773529	335.45	334.8657143	68.11	67.54470588	83.86	83.72617647
273.02	270.1773529	333.81	334.8657143	68.25	67.54470588	83.45	83.72617647
270.7	270.1773529	334.59	334.8657143	67.68	67.54470588	83.65	83.72617647
270.56	270.1773529	336.36	334.8657143	67.64	67.54470588	84.09	83.72617647
272.18	270.1773529	334.89	334.8657143	68.05	67.54470588	83.72	83.72617647
269.9	270.1773529	335.24	334.8657143	67.48	67.54470588	83.81	83.72617647
271.76	270.1773529	334.89	334.8657143	67.94	67.54470588	83.72	83.72617647
270.91	270.1773529	334.65	334.8657143	67.73	67.54470588	83.66	83.72617647
270.31	270.1773529	335.5	334.8657143	67.58	67.54470588	83.88	83.72617647
270.81	270.1773529	331.13	334.8657143	67.7	67.54470588	82.78	83.72617647
271.95	270.1773529	336.46	334.8657143	67.99	67.54470588	84.11	83.72617647
269.19	270.1773529	334.52	334.8657143	67.3	67.54470588	83.63	83.72617647
270.25	270.1773529	335.86	334.8657143	67.56	67.54470588	83.97	83.72617647
271.18	270.1773529	334.62	334.8657143	67.79	67.54470588	83.66	83.72617647
271.26	270.1773529	333.63	334.8657143	67.82	67.54470588	83.41	83.72617647
255.82	270.1773529	335.16	334.8657143	63.96	67.54470588	83.79	83.72617647
271.33	270.1773529	334.69	334.8657143	67.83	67.54470588	83.67	83.72617647
269.35	270.1773529	335.54	334.8657143	67.34	67.54470588	83.89	83.72617647
270.84	270.1773529	334.42	334.8657143	67.71	67.54470588	83.6	83.72617647
270.05	270.1773529	335.81	334.8657143	67.51	67.54470588	83.95	83.72617647
271	270.1773529	337.35	334.8657143	67.75	67.54470588	84.34	83.72617647
270.65	270.1773529	334.55	334.8657143	67.66	67.54470588	83.64	83.72617647
270.86	270.1773529	335.2	334.8657143	67.71	67.54470588	83.8	83.72617647
265.09	270.1773529	335.79	334.8657143	66.27	67.54470588	83.95	83.72617647
267.18	270.1773529	334.45	334.8657143	66.79	67.54470588	83.61	83.72617647

Figure D9. XHPL

Virtual Nodes	Virtual Nodes Avg		Physical Nodes	Physical Nodes Avg	
8.35E-001		8.41E-001	1.80E+000		1.80E+000
8.39E-001		8.41E-001	1.80E+000		1.80E+000
8.38E-001		8.41E-001	1.80E+000		1.80E+000
8.39E-001		8.41E-001	1.80E+000		1.80E+000
8.37E-001		8.41E-001	1.80E+000		1.80E+000
8.39E-001		8.41E-001	1.80E+000		1.80E+000
8.38E-001		8.41E-001	1.80E+000		1.80E+000
8.39E-001		8.41E-001	1.80E+000		1.80E+000
8.37E-001		8.41E-001	1.80E+000		1.80E+000
8.40E-001		8.41E-001	1.80E+000		1.80E+000
8.42E-001		8.41E-001	1.80E+000		1.80E+000
8.39E-001		8.41E-001	1.80E+000		1.80E+000
8.41E-001		8.41E-001	1.80E+000		1.80E+000
8.42E-001		8.41E-001	1.80E+000		1.80E+000
8.40E-001		8.41E-001	1.80E+000		1.80E+000
8.38E-001		8.41E-001	1.80E+000		1.80E+000
8.40E-001		8.41E-001	1.78E+000		1.80E+000
8.41E-001		8.41E-001	1.80E+000		1.80E+000
8.39E-001		8.41E-001	1.80E+000		1.80E+000
8.41E-001		8.41E-001	1.80E+000		1.80E+000
8.41E-001		8.41E-001	1.80E+000		1.80E+000
8.42E-001		8.41E-001	1.78E+000		1.80E+000
8.42E-001		8.41E-001	1.80E+000		1.80E+000
8.40E-001		8.41E-001	1.80E+000		1.80E+000
8.45E-001		8.41E-001	1.80E+000		1.80E+000
8.41E-001		8.41E-001	1.80E+000		1.80E+000
8.43E-001		8.41E-001	1.80E+000		1.80E+000
8.45E-001		8.41E-001	1.80E+000		1.80E+000
8.45E-001		8.41E-001	1.80E+000		1.80E+000
8.43E-001		8.41E-001	1.80E+000		1.80E+000
8.43E-001		8.41E-001	1.80E+000		1.80E+000
8.43E-001		8.41E-001	1.80E+000		1.80E+000
8.42E-001		8.41E-001	1.80E+000		1.80E+000
8.46E-001		8.41E-001	1.80E+000		1.80E+000
8.44E-001		8.41E-001	1.80E+000		1.80E+000